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ΠΣ



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234567

$FF_{\omega}^8$

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<https://5ht.github.io/bertrand/>

[https://llis.nasa.gov/llis\\_lib/pdf/1009464main1\\_0641-mr.pdf](https://llis.nasa.gov/llis_lib/pdf/1009464main1_0641-mr.pdf)

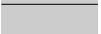
<http://www-users.math.umn.edu/~arnold/disasters/ariane5rep.html>

<http://mdbailey.ece.illinois.edu/publications/imc14-heartbleed.pdf>

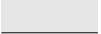
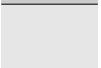
<https://arxiv.org/pdf/1809.03981.pdf>

<https://www.cs.umd.edu/~aseem/solidetherplas.pdf>

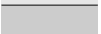
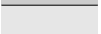
<https://www.dcs.ed.ac.uk/home/mlj/>



$\omega$



$\Pi\Sigma$



\* : \*\* : □□ : □□ : \*P<sub>ω</sub>

$$\begin{array}{ccc} \lambda_2 & \xrightarrow{\lambda_\omega} & \lambda P_\omega \\ \downarrow & & \downarrow \\ \lambda_2 & \xrightarrow{\lambda_\omega} & \lambda P_\omega \\ \downarrow & & \downarrow \\ \lambda_{\rightarrow} & \xrightarrow{\lambda_\omega} & \lambda P_\omega \end{array}$$

[0, 1]

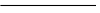


$\mathbf{A}^1$

$\infty$

$\Sigma\Omega$

$\flat\sharp$



1112

1314

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<https://groupoid.github.io/languages>  
<https://arxiv.org/pdf/1804.07608.pdf>  
<https://n2o.dev/ua>  
<https://anders.groupoid.space/lib>





1516

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12

34

567

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<http://mrg.doc.ic.ac.uk/kohei/>

89

Henk  
1011

$\pi\pi$

$O_{CPS}\pi$ Bob  
 $O_{CPS}^{1213}$ Joe

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<https://5ht.co/ctt.pdf>  
<https://5ht.co/cctt.pdf>  
<https://web.sfc.keio.ac.jp/~hagino/thesis.pdf>  
<https://github.com/devaspot/charity>  
<http://nickbenton.name/coqasm.pdf>  
<https://www.cl.cam.ac.uk/~mom22/tphols09-lisp.pdf>

O<sub>λ</sub>λO<sub>π</sub>πO<sub>μ</sub>O<sub>π</sub>O<sub>σ</sub>O<sub>=</sub>O<sub>W</sub>O<sub>I</sub>



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$\pi\Sigma = +\perp TN\mathcal{U}_i E\mathcal{I}$



$\Pi x : \text{op}(o) : \text{Up} : \text{concept}(o)$

$\text{Ob} : \text{UHom} : \text{Ob} \rightarrow \text{Ob} \rightarrow \text{U}$

$\text{id}_o$

$A : \text{Ob}(c) \text{isContr}(\Sigma(B : \text{Ob}(C)), A = B)$

$$p : E \rightarrow Bf : Y \rightarrow Xp(f)Bp(f) = \text{id}_B$$

$$p : E \rightarrow Bf : Y \rightarrow XEt = p(f)u : Z \rightarrow Yp(f) = p(u)a : Z \rightarrow Yf \circ a = u$$

$$p : E \rightarrow Bf : Y \rightarrow XEt = p(f)u : K \rightarrow JBv : Z \rightarrow XEp(v) = t \circ uw : Z \rightarrow Y \\ Ev = f \circ wp(w) = u$$

$$B \rightarrow Ba : \text{Ob}_{\overline{B}} = \text{Hom}_B(x, y)\text{Hom}_{\overline{B}} = [f : \text{Hom}_B, g : \text{Hom}_B]B$$

$$\begin{matrix} f(x) \\ q \\ g(y) \end{matrix}$$

$$p : E \rightarrow BBu : J \rightarrow IBX \in p(I)Bf : Y \rightarrow XuXu$$

$$p^{-1}B^{op} \rightarrow \text{CatPsh}(B) = [B^{op}, \text{Cat}]$$

$$p : X \rightarrow Bq : Y \rightarrow BBF : X \rightarrow Yq \circ F = PxXpF(x)q$$

$$p : E \rightarrow Bq : D \rightarrow AB\text{Fib}_B(p, q)\text{Cat}/Bp \rightarrow q(H : E \rightarrow D, K : B \rightarrow A)fp \\ H(f)q$$

$$p : E \rightarrow B \rightarrow \text{cod} \circ p : E \rightarrow Bf \in \text{Epf}B$$

$$\text{Psh}(B)\text{Fib}(B)$$

$$\int : \text{Psh}(B) \xrightarrow{\cong} \text{Fib}(B)$$

$\Pi\Sigma$

$\text{COb}_C \Pi(A, B) \text{Hom}_C(\Pi(A, B), \Pi(A', B')) [f : A \rightarrow A', g(x : A) : B(x) \rightarrow B'(f(x))]$

```
def CwF : U :=  $\Sigma$  (C: precategory) (T: catfunctor C Fam)
  (context: isContext C) (terminal: isTerminal C), isComprehension C T
```

$CF: C^{op} \rightarrow \text{SetCSet}$

$Ct \in CTy, Tm: C^{op} \rightarrow \text{Setp}: Tm \rightarrow Ty$

`def naturalModel : U :=  $\Sigma$  (C : precategory) ( _ : isCategory C)  
(t : terminal C) (Tm : carrier C) (Ty : carrier C)  
(p : hom C VT V),  $\Pi$  (f : homTo C V), hasPullback C (Tm, f, Ty, p)`

C—C—

$P, Q: C^{op} \rightarrow \alpha: Q \rightarrow P \alpha \text{Ob}(C) x: \text{Ob}(C) p_x: D \rightarrow Cy: Q(D)$

$$\begin{array}{ccc} y(D) & \xrightarrow{p} & y(Q) \\ y(C) & \xrightarrow{p} & y(P) \end{array}$$

$F: C \rightarrow D \phi_{Ty}: F! Ty_C \rightarrow Ty_D \phi_{Tm}: F! Tm_C \rightarrow Tm_D$

$$\begin{array}{ccc} F! Tm_C & \xrightarrow{\phi_{Tm}} & Tm_D \\ F! Ty_C & \xrightarrow{\phi_{Ty}} & Ty_D \end{array}$$

$F!: C^{op} \rightarrow D^{op}$



$O_\lambda$	$\lambda$
$O_\pi$	$\pi$
$O_\mu$	
$O_\Pi$	
$O_\Sigma$	
$O_=_$	
$O_W$	
$O_I$	
$O_\triangleright$	$\pi$
$O/$	
$O_H$	
$O_\dashv$	

TmTy

$\beta\eta$

$O_\Pi O_\Sigma O_=_ O_{PTS} O_{MLTT-80} O_{HTS}$

$$O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow O_I.$$

$$O_{PTS}(O_\Pi) \rightarrow O_{MLTT-72}(O_\Pi, O_\Sigma) \rightarrow O_{MLTT-75}(\dots, O_\Sigma, O_=_) \rightarrow \\ \rightarrow O_{MLTT-80}(\dots, O_=_, O_W) \rightarrow O_{HTS}(\dots, O_W, O_I).$$

$$O_{PTS} : O_\Pi \rightarrow \mathcal{U}$$

$$O_{MLTT-72} : O_\Pi \rightarrow O_\Sigma \rightarrow \mathcal{U}$$

$$O_{MLTT-75} : O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow \mathcal{U}$$

$$O_{MLTT-80} : O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow \mathcal{U}$$

$$O_{HTS} : O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow O_I \rightarrow \mathcal{U}$$

$$O_{HTS} = O_{\Pi\Sigma=WI}$$

$$O_{HTS} = O_{\Pi\Sigma=WI} : O_\Pi \rightarrow O_\Sigma \rightarrow O_=_ \rightarrow O_W \rightarrow O_I \rightarrow \mathcal{U}.$$

$$O_{MLTT-80}$$

$\beta\eta\beta\eta$

$$O_\infty : O_{\text{CPS}} \rightarrow O_{\text{PTS}} \rightarrow O_{\text{MLTT-80}} \rightarrow O_{\text{HTS}} \rightarrow \dots$$

$$f : O_x \rightarrow O_y O_x O_y F_O$$

$$\begin{aligned} \mathcal{F}_L \mathcal{F}f : L \rightarrow L', L \supseteq L \mathcal{F}f : L \rightarrow L', L \subseteq L \mathcal{F}f : L \rightarrow L', L' \cong L \mathcal{F}f : L \rightarrow O\mathcal{F} \\ f : L \rightarrow L' L' \mathcal{F}f : L \rightarrow L' L' \end{aligned}$$

$$O_{I^*} \rightarrow O_{\Pi=} O_{\Pi} \rightarrow O_{\Pi\Sigma} O_{\Pi} \rightarrow O_{\Pi\Sigma} O_{\Pi^*} \rightarrow O_{\Pi}$$

Henk =  $U^n$ ,  $\Pi$ .

Frank =  $U^n$ ,  $\Pi$ , Ind.

Errett =  $U^n$ ,  $\Pi$ ,  $\Sigma$ , Prop.

Per =  $U^n$ ,  $\Pi$ ,  $\Sigma$ , 0, 1, 2, W, Prop.

Christine =  $U^n$ ,  $\Pi$ ,  $\Sigma$ , Id, Ind, Prop

Anders =  $U^n$ ,  $V^n$ ,  $\Pi$ ,  $\Sigma$ , 0, 1, 2, W, Path, Prop.

Urs = Anders,  $U_i^{|\alpha|}$ ,  $A \times B$ ,  $G \rightarrow A$ ,  $s$ ,  $b$ ,  $\#$ ,  $\mathcal{J}$ ,  $\bigcirc$ .

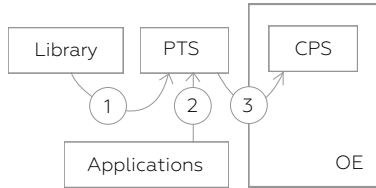
Dan = Anders, Chain, Cochain, Category, Monoid, Group, Ring,  $\Delta$ .

Fabien = Anders,  $k$ ,  $A^1$ ,  $S^{1,1}$ ,  $L_{A^1}$ , Susp, Trunc<sup>n</sup>, Nisn,  $K^1(Z, n)$ , BGL, MGL.

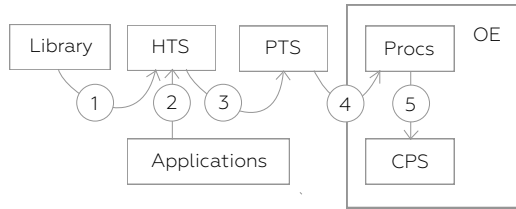
Jack = Anders, Fib<sup>n</sup>, Susp, Trunc<sup>n</sup>,  $\mathbf{N}$ ,  $\mathbf{N}_\infty$ , Spec,  $\pi_n \hat{\sim} S(A)$ ,  $S^o[p]$ , Group,

$A \cup B$ ,  $[A, B]$ ,  $H^n(X; G)$ ,  $G \otimes H$ ,  $SS(E, r)$ .

$$\text{PTS}_{\text{CPS}} = \begin{cases} \text{Ob} : \{\text{O}_{\text{CPS}}, \text{O}_{\text{PTS}}\} \\ \text{Hom} : \{1, 2 : \mathbb{1} \rightarrow \text{O}_{\text{PTS}}, 3 : \text{O}_{\text{PTS}} \rightarrow \text{O}_{\text{CPS}}\} \end{cases}$$



$$\text{Total} = \begin{cases} \text{Ob} : \{O_{\text{CPS}}, O_{\text{PTS}}, O_{\text{MLTT-75}}, O_{\text{MLTT-80}}, O_{\text{HTS}}\} \\ \text{Hom} : \begin{cases} 1, 2 : \mathbb{1} \rightarrow O_{\text{HTS}}, 3 : O_{\text{MLTT-75}} \rightarrow O_{\text{MLTT-80}} \\ 4 : O_{\text{HTS}} \rightarrow O_{\text{MLTT-80}}, 5 : O_{\text{MLTT-80}} \rightarrow O_{\text{PTS}}, 6 : O_{\text{PTS}} \rightarrow O_{\text{CPS}} \end{cases} \end{cases}$$



$O_{\text{CPS}}$

$$O_{\text{CPS}} = \begin{cases} \text{Ob} : \{\text{maybeCPS}\} \\ \text{Hom} : \{\text{eval} : \text{Ob} \rightarrow \text{Ob}\} \end{cases}$$

$O_\lambda O_\pi O_\mu$

$O_{\text{CPS}}$

```
def CPS : U
:= inductive { lambda (c: Joe CPS)
              | process (m: Bob CPS)
              | tensor (f: Alice CPS)
              }
```

$\lambda O_{\text{CPS}} \text{maybe}$

$$O_{\text{CPS}} : O_\lambda \rightarrow O_\pi \rightarrow O_\mu \rightarrow U$$

$O_\lambda$

```
def Joe (cps: U): U
:= inductive { var (x: nat)
              | lam (l: nat) (d: cps)
              | app (f a: cps)
              }
```

$O_\pi$

```
def Bob (lang: U) : U
:= inductive { process (protocol: lang)
              | spawn (cursors: lang) (core: nat) (program: lang)
              | snd (cursor: lang) (data: lang)
              | rcv (cursor: lang)
              | pub (size: nat)
              | sub (cursor: lang)
              }
```

$O_\mu$

```
def Alice (lang: U) : U
:= inductive { Variable (_: Var)
              | Prim (_: Builtin)
              | Star | True | False
              | Int (_: nat) | Float (_: float)
              | Lambda (a: Var) (b: Linear) (c: Exp)
              | App (a b: Exp)
              | Pair (a b: Var) (c d: Exp)
              | Consume (a: Var) (b c: Exp)
              | Gen (a: Var) (b: Exp)
              | Spec (a: Exp) (b: Fraction)
```

```

    | Fix (a b: Var) (c d: Linear) (e: Exp)
    | If (a b c: Exp)
    | Let (a: Var) (b c: Exp)
  }

def Linear : U
:= inductive { Empty | Unit | Bool
             | Int | Float
             | Tensor (a: Fraction) (x: Dimension)
             | Pair (a b: Linear) | Fun (a b: Linear)
             | Consume (a: Linear) | All (a: Var) (b: Linear)
             }

def Builtin : U
:= inductive { Intop (a: Arith) | Floatop (a: Arith)    -- SIMD types
             | Get | Set | Duplicate | Free           -- linearity
             | Transpose | Size                       -- matrices
             | Asum | Axy | Dotp | Rotm | Scal | Amax  -- BLAS Level 1
             | Symm | Gemm | SyrK | Posv              -- BLAS Level 3
             }

def Fraction : U := inductive { Z | S (_: Fraction) }
def Dimension : U := inductive { Vector | Matrix | Stream | Table }
def Arith      : U := inductive { Add | Sub | Mul | Div | Eq | Lt | Gt }

```

$\Pi$   
 $\Sigma$   
 $N \rightarrow u$

$O_{PTS}$

$$O_{PTS} = \begin{cases} \text{Ob} : \{X : \text{maybePTS}, \text{target} : \text{maybeCPS}\} \\ \text{Hom} : \begin{cases} \text{type}, \text{norm} : X \rightarrow X, \text{extract} : X \rightarrow \text{target} \\ \text{certify} : X \rightarrow \text{target} = \text{type} \circ \text{norm} \circ \text{extract} \end{cases} \end{cases}$$

$O_{PTS} O_{PTS} \Pi$

```
def PTS : U := inductive { forall (_, Pi PTS) }  
def Henk := PTS
```

$O_{\Pi}$

```
def Pi (lang: U) : U  
:= inductive { fibrant (n: nat)  
  | variable (x: name) (l: nat)  
  | pi (x: name) (l: nat) (f: lang)  
  | lambda (x: name) (l: nat) (f: lang)  
  | application (f a: lang)  
}
```

$\Pi\Sigma\text{MLTT} - 75\Pi\Sigma^2\Pi\nabla^3$

$O_{\text{MLTT}-75}$

$$O_{\text{MLTT}-75} = \begin{cases} \text{Ob} : \{\text{maybeMLTT} - 75\} \\ \text{Hom} : \begin{cases} \text{type, norm} : \text{Ob} \rightarrow \text{Ob} \\ \text{certify} : \text{Ob} \rightarrow \text{Ob} = \text{type} \circ \text{norm} \end{cases} \end{cases}$$

$O_{\text{MLTT}-75}O_{\text{MLTT}-75}O_{\Pi}O_{\Sigma}O_{=}$

```
def MLTT : U
:= inductive { forall (_, Pi MLTT)
  | sigma (_, Sigma MLTT)
  | id (_, Id MLTT)
}
```

$O_{\Sigma}O_{\text{MLTT}-72}O_{\Pi\Sigma}$

```
def Sigma (lang: U) : U
:= inductive { sigma (n: name) (a b: lang)
  | pair (a b: lang)
  | fst (p: lang)
  | snd (p: lang)
}
```

$O_{=}O_{\text{MLTT}-75}O_{\Pi\Sigma=}$

```
def Id (lang: U) : U
:= inductive { identity (t a b: lang)
  | id_intro (a b: lang)
  | id_elim (a b c d e: lang)
  | id_compute (a b c d e: lang)
}
```

$O_{=\eta}$

$O_{\text{MLTT-80}}$

$$O_{\text{MLTT-80}} = \begin{cases} \text{Ob} : \{X : \text{maybePM}, \text{target} : \text{maybeCPS}\} \\ \text{Hom} : \begin{cases} \text{type, norm, induction} : X \rightarrow X, \text{extract} : X \rightarrow \text{target} \\ \text{certify} : X \rightarrow \text{target} \\ \text{certify} = \text{type} \circ \text{norm} \circ \text{induction} \circ \text{extract} \end{cases} \end{cases}$$

$O = O_{\Sigma} O_{\Pi} O_0 O_1 O_2 O_W$

$O_{\text{MLTT-80}}$

```
def MLTT-80 := Per
```

```
def Per : U
```

```
:= inductive { forall (_: Pi Per)
             | sigma (_: Sigma Per)
             | id (_: Id Per)
             | 0 (_: Empty Per)
             | 1 (_: Unit Per)
             | 2 (_: Bool Per)
             | W (_: W Per)
             }
```

```
def W (lang: U) : U
```

```
:= inductive { W_Form (n: name) (a b: lang)
             | W_Sup (a b: lang)
             | W_Ind (a b c: lang)
             }
```

```
def Empty (lang: U) : U := inductive { 0_Ind (a: lang) }
```

```
def Unit (lang: U) : U := inductive { unit | star | 1_Ind (a: lang) }
```

```
def Bool (lang: U) : U := inductive { bool | true | false | 2_Ind (a: lang)
                                     }
```

$O_{CIC}$

$$O_{CIC} = \begin{cases} \text{Ob} : \{X : \text{maybePM}, \text{target} : \text{maybeCPS}\} \\ \text{Hom} : \begin{cases} \text{type}, \text{norm}, \text{induction} : X \rightarrow X, \text{extract} : X \rightarrow \text{target} \\ \text{certify} : X \rightarrow \text{target} \\ \text{cerfity} = \text{type} \circ \text{norm} \circ \text{induction} \circ \text{extract} \end{cases} \end{cases}$$

$O = O_{\Sigma} O_{\Pi}$

$O_{CIC}$

```
def Frank := inductive { forall (_, Pi Frank) | ind (_, Ind Frank) }
def Christine := CIC
def CIC : U
:= inductive { forall (_, Pi CIC)
  | sigma (_, Sigma CIC)
  | id (_, Id CIC)
  | prop (_, Id CIC)
  | ind (_, Ind CIC)
}
```

$O_{IND} O_{IND}$

```
def Ind (lang: U) : U
:= inductive { formation (_, Inductive lang)
  | constructor (_, list (triple Nat Inductive lang))
  | eliminator (t: Inductive) (a b: lang) (cases: list lang)
}

def Inductive (lang: U) : U
:=  $\Sigma$  (name : string)
  (params : list (prod name lang))
  (level : Nat)
  (constrs : list (prod (Nat lang))), 1
```

$O_{HTS}$

$$O_{HTS} = \begin{cases} \text{Ob} : \{\text{maybeHTS}\} \\ \text{Hom} : \begin{cases} \text{type, norm} : \text{Ob} \rightarrow \text{Ob} \\ \text{certify} : \text{Ob} \rightarrow \text{Ob} = \text{type} \circ \text{norm} \end{cases} \end{cases}$$

$O_{HTS} O_I O_W O = O_\Sigma O_\Pi$

```
def HTS : U
:= inductive { forall (_, Pi HTS)
  | sigma (_, Sigma HTS)
  | id (_, Id HTS)
  | prop (_, Id HTS)
  | 0 (_, Empty HTS)
  | 1 (_, Unit HTS)
  | 2 (_, Bool HTS)
  | W (_, W HTS)
  | homotopy (_, Homotopy HTS)
}
```

$O_{MLTT-80}$

$O_I$

```
def CCHM (lang: U) : U
:= inductive { pretype (n: nat)
  | PathP (_, lang) | PLam (_, lang) | PApp (f a: lang)
  | I | 0 | 1 | And (a b: lang) | Or (a b: lang) | Neg (_, lang)
  | Transp (a b: lang) | HComp (a b c d: lang)
  | Partial (_, lang) | PartialP (a b: lang) | System (_, lang)
  | Sub (a b c: lang) | Inc (a b: lang) | Ouc (: lang)
  | Glue (: lang) | GlueElem (a b c: lang) | Unglue (_, lang)
}
```

$O_{HTS} O = O_I$



```

fun a (0, n) = n + 1
  | a (m, 0) = a (m - 1, 1)
  | a (m, n) = a (m - 1, a (m, n - 1))

```

```

fun proc =
let val p0 = pub(0,8)
    val s1 = sub(0,p0)
    val s2 = sub(0,p0)
in send(p0,11);
  send(p0,12);
  [ receive(s1);
    receive(s2);
    receive(s1);
    receive(s2)
  ]
end

```

```

fun simpleConvolution (i n: int) (x0: float) (write w: vector float)
: vector float
= begin
  if n = i then result.emit(write),
  a = [w0,w1,w2] = w.get(0,3),
  b = [x0,x1,x2] = [ x0 | write.get(i,2) ],
  write.set(i, Dotp(a,b)),
  simpleConvolution((i + 1),n,x1,write,w)
end

```

$\Pi\Sigma$

```
def zero : (T → T) → T → T := λ (s: T → T) (z: T), z
def succ : ((T → T) → T → T) → ((T → T) → T → T)
:= λ (w: (T → T) → T → T) (y: T → T) (x: T), y (w y x)
```

$\Pi$

```
def N := Π (A : U), (A → A) → A → A
def zero : N := λ (A : U) (S : A → A) (Z : A), Z
def succ : N → N := λ (n : N) (A : U) (S : A → A) (Z : A), S (n A S Z)
def plus (m n : N) : N := λ (A : U) (S : A → A) (Z : A), m A S (n A S Z)
def mult (m n : N) : N := λ (A : U) (S : A → A) (Z : A), m A (n A S) Z
def pow (m n : N) : N := λ (A : U) (S : A → A) (Z : A), n (A → A) (m A) S Z
```

$\Pi\Sigma$

```
def empty      : U := inductive { }
def L1 (A : U) : U := inductive { nil | cons (head: A) (tail: L1 A) }
def S1        : U := inductive { base | loop : Equ S1 base base }

def quot (A: U) (R : A → A → U) : U
:= inductive { quotient (a: A)
              | identification (a b: A) (r: R a b)
              : Equ (quot A R) (quotient a) (quotient b)
              }
```

```

def idfun (A : U) : A → A := (λ (a : A), a)
def idfun' (A : U) : A → A := transp (<i>A) 0
def idfun'' (A : U) : A → A := (λ (a : A), hcomp A 0 (λ (i : I), []) a)
def isFiberBundle (B: U) (p: B → U) (F: U): U
  := ∑ (v: U) (w: surjective v B), (Π (x: v), PathP (<_>U) (p (w.1 x)) F)
def ~- (X : U) (a x' : X) : U := Path (∫ X) (ι X a) (ι X x')
def ID (X : U) (a : X) : U := ∑ (x' : X), ~- X a x'
def unitDisc (X : U) (x : ∫ X) : U := ∑ (x' : X), Path (∫ X) x (ι X x')
def starDisc (X : U) (x : X) : ID X x := (x, idp (∫ X) (ι X x))
def T∞ (A : U) : U := ∑ (a : A), ID A a
def inf-prox-ap (X Y : U) (f : X → Y) (x x' : X) (p : ~- X x x')
  : ~- Y (f x) (f x') := <i>∫-app X Y f (p @ i)
def d (X Y : U) (f : X → Y) (x : X) (e : ID X x) : ID Y (f x)
  := (f e.1, inf-prox-ap X Y f x e.1 e.2)
def T∞-map (X Y : U) (f : X → Y) (τ : T∞ X) : T∞ Y
  := (f τ.1, d X Y f τ.1 τ.2)
def is-homogeneous (A : U)
  := ∑ (e : A) (t : A → equiv A A),
    Π (x : A), Path A ((t x).1 e) x

```






$O_{CPS_\lambda}$

```
def CPS $\lambda$  : U
:= inductive { var (x: nat)
              | lam (l: nat) (d: cps)
              | app (f a: cps)
              }
```

$O_{CPS}$





```
objdump ./target/release/o -d | grep mulpd
223f1: c5 f5 59 0c d3    vmulpd (%rbx,%rdx,8),%ymm1,%ymm1
223f6: c5 dd 59 64 d3 20 vmulpd 0x20(%rbx,%rdx,8),%ymm4,%ymm4
22416: c5 f5 59 4c d3 40 vmulpd 0x40(%rbx,%rdx,8),%ymm1,%ymm1
2241c: c5 dd 59 64 d3 60 vmulpd 0x60(%rbx,%rdx,8),%ymm4,%ymm4
2264d: c5 f5 59 0c d3    vmulpd (%rbx,%rdx,8),%ymm1,%ymm1
22652: c5 e5 59 5c d3 20 vmulpd 0x20(%rbx,%rdx,8),%ymm3,%ymm3
```

## O\_CPS

```
def Lazy : U
:= inductive { Defer (otree: NodeId) (a: AST) (cont: Cont)
| Continuation (otree: NodeId) (a: AST) (cont: Cont)
| Return (a: AST)
| Start
}

def Cont : U
:= inductive { Expressions (a: AST) (v: Option (Iter AST)) (c: Cont)
| Assign (ast: AST) (cont: Cont)
| Cond (c,d: AST) (cont: Cont)
| Func (a,b,c: AST) (cont: Cont)
| List (acc: Vec AST) (vec: Iter AST) (i: Nat) (c: Cont)
| Call (a: AST) (i: Nat) (cont: Cont)
| Return
| Intercore (m: Message) (cont: Cont)
| Yield (cont: Cont)
}
```

## O\_CPS

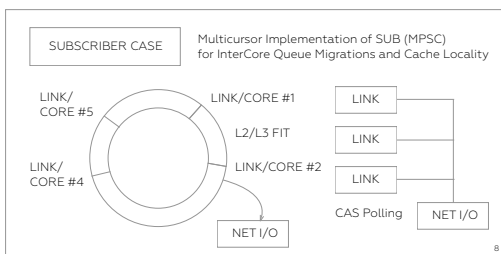
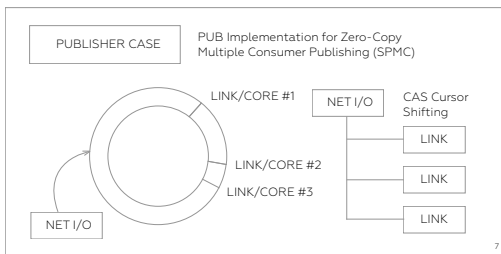
```
E: V | A | C
NC: ";" = [] | ";" m:NL = m
FC: ";" = [] | ";" m:FL = m
EC: ";" = [] | ";" m:EL = m
NL: NAME | o:NAME m:NC = Cons o m
FL: E | o:E | m:FC = Cons o m
EL: E | EC | o:E m:EC = Cons o m
C: N | c:N a:C = Call c a
N: NAME | S | HEX | L | F
L: "(" ")" = [] | "(" [" c:NL "]" m:FL ")" = Table c m
    | "(" " 1:EL ")" = List l
F: "{" "}" = Lambda [] [] []
    | "{" [" c:NL "]" m:EL "}" = Lambda [] c m
    | "{" m:EL "}" = Lambda [] [] m
```

```
def AST : U
:= inductive { Atom (a: Scalar)
              | Vector (a: Vec AST)
            }
```

```
def Value : U
:= inductive { Nil
              | SymbolInt (a: u16)
              | SequenceInt (a: u16)
              | Number (a: i64)
              | Float (a: f64)
              | VecNumber (Vec i64)
              | VecFloat (Vec f64)
            }
```

```
def Scalar : U
:= inductive { Nil
| Any
| List (a: AST)
| Dict (a: AST)
| Call (a b: AST)
| Assign (a b: AST)
| Cond (a b c: AST)
| Lambda (otree: Option NodeId) (a b: AST)
| Yield (c: Context)
| Value (v: Value)
| Name (s: String)
}
```

```
reactor[aux;0;mod[console;network]];
reactor[timercore;1;mod[timer]];
reactor[core1;2;mod[task]];
reactor[core2;3;mod[task]];
```





```
fun pub(capacity: int): int
```

```
fun sub(publisher: int): int
```

```
fun spawn(core: int, program: code, cursors: array int): int
```

```
fun kill(process: int): int
```

```
fun send(writer: int, data: binary): int
```

```
fun receive(reader: Int)
```

## O<sub>CPS</sub>

```
pub struct Publisher<T> {  
    ring: Arc<RingBuffer<T>>,  
    next: Cell<Sequence>,  
    cursors: UncheckedUnsafeArc<Vec<Cursor>>,  
}
```

```
pub struct Subscriber<T> {  
    ring: Arc<RingBuffer<T>>,  
    token: usize,  
    next: Cell<Sequence>,  
    cursors: UncheckedUnsafeArc<Vec<Cursor>>,  
}
```

```
pub struct Channel {  
    publisher: Publisher<Message>,  
    subscribers: Vec<Subscriber<Message>>,  
}
```

```
pub struct Memory<'a> {  
    publishers: Vec<Publisher<Value<'a>>>,  
    subscribers: Vec<Subscriber<Value<'a>>>,  
}
```

```
pub struct Scheduler<'a> {  
    pub tasks: Vec<T3<Job<'a>>>,  
    pub bus: Channel,  
    pub queues: Memory<'a>,  
    pub io: IO,  
}
```

```
pub enum Message {
  Pub(Pub),
  Sub(Sub),
  Print(String),
  Spawn(Spawn),
  AckSub(AckSub),
  AckPub(AckPub),
  AckSpawn(AckSpawn),
  Exec(usize, String),
  Select(String, u16),
  QoS(u8, u8, u8),
  Halt,
  Nop,
}
```

```
def AVX-512 : U
:= inductive { Star | True | False
  | Variable (_: Var)
  | Prim (_: Builtin)
  | Int (_: nat) | Float (_: float)
  | Lambda (a: Var) (b: Linear) (c: Exp)
  | App (a b: Exp)
  | Pair (a b: Var) (c d: Exp)
  | Consume (a: Var) (b c: Exp)
  | Gen (a: Var) (b: Exp)
  | Spec (a: Exp) (b: Fraction)
  | Fix (a b: Var) (c d: Linear) (e: Exp)
  | If (a b c: Exp)
  | Let (a: Var) (b c: Exp)
}
```



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<https://github.com/o1>  
<https://github.com/o3>  
<https://ws.erp.uno>  
<https://github.com/o89>



5

6



${}^7F_{\omega}$

```
axiom all : IO (List TableId)
axiom browse : IO String
axiom delete :  $\Pi$  (k: U), TableId -> k -> IO Unit
axiom first :  $\Pi$  (t: TableId) (k: U), IO (Maybe k)
axiom last :  $\Pi$  (k: U), TableId -> IO (Maybe k)
axiom foldl :  $\Pi$  (v acc: U), (v -> acc -> acc) -> acc -> TableId -> IO acc
axiom foldr :  $\Pi$  (v acc: U), (v -> acc -> acc) -> acc -> TableId -> IO acc
axiom insert :  $\Pi$  (v: U), TableId -> v -> IO Boolean
axiom lookup :  $\Pi$  (k v: U), TableId -> k -> IO (List v)
axiom member :  $\Pi$  (k: U), TableId -> k -> IO Boolean
axiom new : Atom -> TableOptions -> IO TableId
axiom next :  $\Pi$  (k: U), TableId -> k -> IO (Maybe k)
axiom prev :  $\Pi$  (k: U), TableId -> k -> IO (Maybe k)
axiom rename : TableId -> Atom -> IO Atom
axiom take :  $\Pi$  (k v: U), TableId -> k -> IO (List v)
axiom match :  $\Pi$  (a v: U), TableId -> a -> IO (List v)
axiom slot :  $\Pi$  (v: U), TableId -> Integer -> IO (Maybe (List v))
```

```
axiom pickle : Binary -> Binary
axiom depickle : Binary -> Binary
axiom encode :  $\prod (k: U), k \rightarrow \text{Binary}$ 
axiom decode :  $\prod (k: U), \text{Binary} \rightarrow \text{IO } k$ 
axiom reg:  $\prod (k: U), k \rightarrow \text{IO } k$ 
axiom unreg :  $\prod (k: U), k \rightarrow \text{IO } k$ 
axiom send :  $\prod (k v z: U), k \rightarrow v \rightarrow \text{IO } z$ 
axiom getSession :  $\prod (k v: U), k \rightarrow \text{IO } v$ 
axiom putSession :  $\prod (k v: U), k \rightarrow v \rightarrow \text{IO } v$ 
axiom getCache :  $\prod (k v: U), \text{Atom} \rightarrow k \rightarrow \text{IO } v$ 
axiom putCache :  $\prod (k v: U), \text{Atom} \rightarrow k \rightarrow v \rightarrow \text{IO } v$ 
```

```
data PI = PI String Atom Atom Atom Integer RestartType
data Sup = Ok Pid String | Error String
```

```
axiom start : PI -> IO Sup
```

```

axiom get :  $\Pi$  (f k v: U), f -> k -> IO (Maybe v)
axiom put :  $\Pi$  (r: U), r -> IO StoreResult
axiom delete :  $\Pi$  (f k: U), f -> k -> StoreResult
axiom index :  $\Pi$  (f p v r: U), f -> Atom -> v -> List r

data Reader = Reader Integer Binary ETF String Integer
data Writer = Writer Integer Binary ETF String Integer
data StoreResult = Ok Integer String Binary
                  | Error Integer String Binary

axiom next : Reader -> IO Reader
axiom prev : Reader -> IO Reader
axiom take : Reader -> IO Reader
axiom drop : Reader -> IO Reader
axiom save : Reader -> IO Reader
axiom append :  $\Pi$  (f r: U), f -> r -> IO StoreResult
axiom remove :  $\Pi$  (f r: U), f -> r -> IO StoreResult

axiom start : Proc -> IO Sup
axiom stop : String -> IO Sup
axiom next : ProcId -> IO ProcRes
axiom load : ProcId -> IO ProcRes
axiom proc : ProcId -> IO ProcRes
axiom assign : ProcId -> IO ProcRes
axiom persist : ProcId -> Proc -> IO ProcRes
axiom amend :  $\Pi$  (k: U), ProcId -> k -> IO ProcRes
axiom discard :  $\Pi$  (k: U), ProcId -> k -> IO ProcRes
axiom modify :  $\Pi$  (k: U), ProcId -> k -> Atom -> IO ProcRes
axiom event : ProcId -> String -> IO ProcRes
axiom head : ProcId -> IO Hist
axiom hist : ProcId -> IO (List Hist)
axiom step : ProcId -> Atom -> IO (List Hist)
axiom docs : ProcId -> IO (List Tuple)
axiom events : ProcId -> IO (List Tuple)
axiom tasks : ProcId -> IO (List Tuple)
axiom doc : Tuple -> ProcId -> IO (List Tuple)

data ProcId = String
data Proc = Proc ProcId String
data ProcRes = Ok Integer String Binary
              | Error Integer String Binary

axiom q :  $\Pi$  (k: U), Atom -> k
axiom qc :  $\Pi$  (k: U), Atom -> k

```

```

axiom jse : Maybe Binary -> Binary
axiom hte : Maybe Binary -> Binary
axiom wire (actions: List Action) : IO (List Action)
axiom render (content: Either Action Element) : Binary
axiom insert_top (dom: Atom) (content: List Element) : IO (List Action)
axiom insert_bottom (dom: Atom) (content: List Element) : IO (List Action)
axiom update (dom: Atom) (content: List Element) : IO (List Action)
axiom clear (dom: Atom) : IO Unit
axiom remove (dom: Atom) : IO Unit

```

```

#textbox { id=userName, body= <<"Anonymous">> },
#panel { id=chatHistory, class=chat_history }

```

```

<input value="Anonymous" id="userName" type="text"/>
<div id="chatHistory" class="chat_history"></div>

```

```

nitro:update(loginButton,
  #button{id = loginButton,
    body = "Login",
    postback = login,
    source = [user, pass]});

```

```

qi('loginButton').outerHTML='<button id=\"loginButton\"
type=\"button\">Login</button>'; { var x=qi('loginButton');
x && x.addEventListener('click',function (event){
  event.preventDefault(); { if (validateSources(['user','pass'])) {
    ws.send(enc(tuple(atom('pickle'),bin('loginButton'),
      bin('b840bca20b3295619d1157105e355880f850bf0223f2f081549dc
8934ecbcd3653f617bd96cc9b36b2e7a19d2d47fb8f9fbe32d3c768866
cb9d6d85700416edf47b9b90742b0632c750a4240a62dfc56789e0f5d8
590f9afdfb31f35fbab1563ec54fcb17a8e3bad463218d6402f1304'),
      [tuple(tuple(string('loginButton'),bin('detail')),[]),
        tuple(atom('user'),querySource('user')),
        tuple(atom('pass'),querySource('pass'))]])); }
    else console.log('Validation Error'); });});

```

8

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$123 \prod \Sigma^\infty$

PTSCoCHenkAlonzo

4

$\omega$

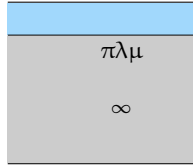
Morte<sup>5</sup>cubicaltt<sup>6</sup>

$$\left\{ \begin{array}{l} \text{Sorts} = \mathbb{U}. \{i\}, i : \text{Nat} \\ \text{Axioms} = \mathbb{U}. \{i\} : \mathbb{U}. \{\text{inci}\} \\ \text{Rules} = \mathbb{U}. \{i\} \sim \mathbb{U}. \{j\} : \mathbb{U}. \{\text{maxij}\} \end{array} \right.$$

HenkHenk

Henk $\Pi\lambda\beta\eta$ Henk

---



7

ω

ω

ω

```

I := list nat
U := U + Unat
0 := U + I + ( 0 ) + 0 0 + 0 → 0
  + λ ( I : 0 ) → 0
  + ∀ ( I : 0 ) → 0

```

→λ∀O<sub>PTS</sub>O<sub>Π</sub>

```

def PTS (lang: U) : U
:= inductive { star      (n: nat)
  | var    (x: name) (l: nat)
  | pi     (x: name) (l: nat) (d c: lang)
  | lambda (x: name) (l: nat) (d c: lang)
  | app    (f a: lang)
}

```

∞<sub>ω</sub><sup>8</sup> Fixpoint

U<sub>0</sub> : U<sub>1</sub> : U<sub>2</sub> : U<sub>3</sub> : ...

U<sub>0</sub>U<sub>1</sub>U<sub>2</sub>U<sub>3</sub>

$\frac{o : \text{Nat}}{U_o}$

U : U

$$A_1 \quad \frac{i : \text{Nat}, j : \text{Nat}, i < j}{U_i : U_j}$$

$$R_1 \quad \frac{i : \text{Nat}, j : \text{Nat}}{U_i \rightarrow U_j : U_{\max(i,j)}}$$

$$A_2 \quad \frac{i : \text{Nat}}{U_i : U_{i+1}}$$

$$R_2 \quad \frac{i : \text{Nat}, j : \text{Nat}}{U_i \rightarrow U_j : U_j}$$

list Sigma

$$1 \quad \frac{\overline{\Gamma : \text{Ctx}}}{\frac{\Gamma : \text{Ctx}}{\emptyset : \Gamma}}$$

$$2 \quad \frac{A : U_i, x : A, \Gamma : \text{Ctx}}{(x : A) \vdash \Gamma : \text{Ctx}}$$

$\beta\eta$ O<sub>P</sub>T<sub>S</sub> $\beta\eta$

$$\Pi \quad \frac{A : U_i, x : A \vdash B : U_j}{\Pi(x : A) \rightarrow B : U_{p(i,j)}}$$

$$\lambda \quad \frac{x : A \vdash b : B}{\lambda(x : A) \rightarrow b : \Pi(x : A) \rightarrow B}$$

$$\text{App} \quad \frac{f : (\Pi(x : A) \rightarrow B)a : A}{fa : B[a/x]}$$

$$\beta \quad \frac{x : A \vdash b : Ba : A}{(\lambda(x : A) \rightarrow b)a = b[a/x] : B[a/x]}$$

$$\frac{\pi_1 : Au : A \vdash \pi_2 : B}{[\pi_1/u]\pi_2 : B}$$

O=O<sub>I</sub>

remoteremote

```
type (:star,N)    D → (:star,N+1)
(:var,N,I)       D → :true = proplists:is_defined N B, om:keyget N D I
(:pi,N,0,I,0)   D → (:star,h(star(type I D)),star(type 0 [(N,norm I)|D]))
(:fn,N,0,I,0)   D → let star (type I D), NI = norm I
                  in (:pi,N,0,NI,type(0,[(N,NI)|D]))
(:app,F,A)      D → let T = type(F,D),
                  (:pi,N,0,I,0) = T, :true = eq I (type A D)
                  in norm (subst 0 N A)
```

```
sh (:star,X)     N P → (:star,X)
(:var,N,I)       N P → (:var,N,I+1) when I >= P
                  → (:var,N,I)
(:pi,N,0,I,0)   N P → (:pi,N,0,sh I N P,sh 0 N P+1)
(:fn,N,0,I,0)   N P → (:fn,N,0,sh I N P,sh 0 N P+1)
(:app,L,R)      N P → (:app,L,R)
```

```

sub (:star,X)      N V L → (:star,X)
  (:var,N,L)      N V L → V
  (:var,N,I)      N V L → (:var,N,I-1) when I > L
  (:pi,N,0,I,0)  N V L → (:pi,N,0,sub I N V L,sub 0 N (sh V N 0) L+1)
  (:pi,F,X,I,0)  N V L → (:pi,F,X,sub I N V L,sub 0 N (sh V F 0) L)
  (:fn,N,0,I,0)  N V L → (:fn,N,0,sub I N V L,sub 0 N (sh V N 0) L+1)
  (:fn,F,X,I,0)  N V L → (:fn,F,X,sub I N V L,sub 0 N (sh V F 0) L)
  (:app,F,A)     N V L → (:app, sub F N V L,sub A N V L)

```

```

norm (:star,X)    → (:star,X)
  (:var,X)        → (:var,X)
  (:pi,N,0,I,0)  → (:pi,N,0,norm I,norm 0)
  (:fn,N,0,I,0)  → (:fn,N,0,norm I,norm 0)
  (:app,F,A)     → case norm F of
                    (:fn,N,0,I,0) → norm (subst 0 N A)
                    NF → (:app,NF,norm A) end

```

```

eq (:star,N)      (:star,N)      → true
  (:var,N,I)      (:var,(N,I))    → true
  (:pi,N1,0,I1,01) (:pi,N2,0,I2,02) →
    let :true = eq I1 I2
    in eq 01 (subst (shift 02 N1 0) N2 (:var,N1,0) 0)
  (:fn,N1,0,I1,01) (:fn,N2,0,I2,02) →
    let :true = eq I1 I2
    in eq 01 (subst (shift 02 N1 0) N2 (:var,N1,0) 0)
  (:app,F1,A1)    (:app,F2,A2)    → let :true = eq F1 F2 in eq A1 A2
  (A,B)           → (:error,(:eq,A,B))

```

```
> ./om help me
[{a,[expr],"to parse. Returns {_,_} or {error,_}."},
 {type,[term],"typechecks and returns type."},
 {erase,[term],"to untyped term. Returns {_,_}."},
 {norm,[term],"normalize term. Returns term's normal form."},
 {file,[name],"load file as binary."},
 {str,[binary],"lexical tokenizer."},
 {parse,[tokens],"parse given tokens into {_,_} term."},
 {fst,[{x,y}], "returns first element of a pair."},
 {snd,[{x,y}], "returns second element of a pair."},
 {debug,[bool],"enable/disable debug output."},
 {mode,[name],"select metaverse folder."},
 {modes,[],"list all metaverses."}]

> ./om print fst erase norm a "#List/Cons"
  \ Head
-> \ Tail
-> \ Cons
-> \ Nil
-> Cons Head (Tail Cons Nil)
ok
```

$(\mu_{L_A}, \text{in})_{L_A}(X) = 1 + (A \times X)\mu_{L_A} = \text{List}(A)\text{nil} : 1 \rightarrow \text{List}(A)\text{cons} : A \times \text{List}(A) \rightarrow \text{List}(A)\text{nil} = \text{in} \circ \text{inl}\text{cons} = \text{in} \circ \text{inr}\text{in} = [\text{nil}, \text{cons}]$   
 $c : 1 \rightarrow \text{Ch} : A \times C \rightarrow C f = ([c, h]) : \text{List}(A) \rightarrow C$

$$\begin{cases} f \circ \text{nil} = c \\ f \circ \text{cons} = h \circ (\text{id} \times f) \end{cases}$$

$f = \text{foldr}(c, h)\mu(1 + A \times X)[1 \rightarrow \text{List}(A), A \times \text{List}(A) \rightarrow \text{List}(A)]$   
 $\text{OPTS}$

$$\begin{cases} \text{foldr} = ([f \circ \text{nil}, h]), f \circ \text{cons} = h \circ (\text{id} \times f) \\ \text{len} = ([\text{zero}, \lambda n \rightarrow \text{succ}n]) \\ (++) = \lambda x \text{sys} \rightarrow ([\lambda(x) \rightarrow \text{ys}, \text{cons}]) (\text{xs}) \\ \text{map} = \lambda f \rightarrow ([\text{nil}, \text{cons} \circ (f \times \text{id})]) \end{cases}$$

```

def list (A: U) : U
:= inductive { cons (x: A) (cs: list A)
             | nil
             }

```

$$\begin{cases} \text{list} = \lambda \text{ctor} \rightarrow \lambda \text{cons} \rightarrow \lambda \text{nil} \rightarrow \text{ctor} \\ \text{cons} = \lambda x \rightarrow \lambda \text{xs} \rightarrow \lambda \text{list} \rightarrow \lambda \text{cons} \rightarrow \lambda \text{nil} \rightarrow \text{cons}x(\text{xs}\text{list}\text{cons}\text{nil}) \\ \text{nil} = \lambda \text{list} \rightarrow \lambda \text{cons} \rightarrow \lambda \text{nil} \rightarrow \text{nil} \end{cases}$$

```

axiom map (A B: U) (f: A -> B) : list A -> list B
axiom length (A: U): list A -> nat
axiom append (A: U): list A -> list A -> list A
axiom foldl (A B: U) (f: B -> A -> B) (Z: B): list A -> B
axiom filter (A: U) (p: A -> bool) : list A -> list A

```

$$\begin{cases} \text{len} = \text{foldr}(\lambda x n \rightarrow \text{succ}n)0 \\ (++) = \lambda \text{ys} \rightarrow \text{foldr}\text{cons}\text{ys} \\ \text{map} = \lambda f \rightarrow \text{foldr}(\lambda x \text{xs} \rightarrow \text{cons}(f x)\text{xs})\text{nil} \\ \text{filter} = \lambda p \rightarrow \text{foldr}(\lambda x \text{xs} \rightarrow \text{if } p x \text{ then } \text{cons } x \text{xs} \text{ else } \text{xs})\text{nil} \\ \text{foldl} = \lambda f v \text{xs} = \text{foldr}(\lambda x g \rightarrow (\lambda \rightarrow g(fax)))\text{id } x \text{sv} \end{cases}$$

$$F_A(X) = 1 + A \times X \quad F_A(X) = A + X \times X \quad 1 + F_A(X) = A \times X$$

$$\mu X \rightarrow 1 + X$$

$$\mu X \rightarrow 1 + A \times X$$

$$\mu X \rightarrow 1 + X \times X + X$$

$$\nu X \rightarrow A \times X$$

$$\nu X \rightarrow 1 + A \times X$$

$$\mu X \rightarrow \mu Y \rightarrow 1 + X \times Y = \mu X = \text{List } X$$

ΠΠΠ [FixpointfunExthomotopy](#)

## Henkλλ

```
ext (:var,X,N,F)      → (:var,X)
  (:app,A,B,N,F)     → (:call,N,ext(F,A,N),[ext(F,B,N)])
  (:fn,S,_,I,O,N,F)  → (:fun,N,(:clauses,[{:clause,N,
                                         [(:var,N,S)],[],[ext(F,O,N)]}]))
  _ → []
```

∞Henk

HenkOPTSHenkHenk

W012

$$W \quad \frac{A : \text{Type} \quad x : A \quad B(x) : \text{Type}}{W(x : A) \rightarrow B(x) : \text{Type}}$$

$$W \quad \frac{a : A \quad t : B(a) \rightarrow W}{\text{sup}(a, t) : W}$$

$$W \quad \frac{\begin{array}{l} w : W \vdash C(w) : \text{Type} \\ x : A, u : B(x) \rightarrow W, \\ v : \Pi(y : B(x)) \rightarrow C(u(y)) \vdash c(x, u, v) : C(\text{sup}(x, u)) \end{array}}{w : W \vdash \text{wrec}(w, c) : C(w)}$$

$$W\beta \quad \frac{\begin{array}{l} w : W \vdash C(w) : \\ x : A, u : B(x) \rightarrow W, \\ v : \Pi(y : B(x)) \rightarrow C(u(y)) \vdash c(x, u, v) : C(\text{sup}(x, u)) \end{array}}{x : A, u : B(x) \rightarrow W \vdash \text{wrec}(\text{sup}(x, u), c) \\ = c(x, u, \lambda(y : B(x)), \text{wrec}(u(y), c)) : C(\text{sup}(x, u))}$$

WMLTTPer

## CICFrankChristine

```
file -> 'module' id 'do' content : {module,'2','4'} .
content -> 'empty' : [] .
content -> line content : ['1'|'2'] .
line -> 'import' id : {import,'2'} .
line -> 'def' id tele ':' exp ':' exp : {def,'2','5','3','7'} .
tele -> 'empty' : [] .
tele -> optic tele: ['1'|'2'] .
optic -> '(' vars ':' exp ')' : {'()', '2', '4'} .
vars -> id : '1' .
vars -> id vars : ['1'|'2'] .
exp -> id : '1' .
exp -> id '.' exp : {'.', '1', '3'} .
exp -> exp exp : {app, '1', '2'} .
exp -> '(' exp ')' : '2' .
exp -> forall tele ',' exp : {'Π', '2', '4'} .
exp -> sigma tele ',' exp : {'Σ', '2', '4'} .
exp -> exp arrow exp : →{'', '1', '3'} .
exp -> lam tele ',' exp : {'λ', '2', '4'} .
exp -> '?' : {hole} .
exp -> U : {'U'} .
exp -> app : '1' .
exp -> 'inductive' '{ sum }' : {inductive, lists:flatten(uncons('3'))} .
constructor -> id tele : {element, '1', '2'} .
constructor -> id tele ':' exp : {interval, '1', '2', '4'} .
sum -> 'empty' : [] .
sum -> constructor : '1' .
sum -> constructor '|' sum : ['3'|'1'] .
app -> exp exp : {app, '1', '2'} .
```

$I : \square_n^{\text{op}} \rightarrow \text{SetCCHMHTSAnders}$

```

start <Module.file> file
start <Module.command> repl
repl : COLON IDENT exp1 EOF | COLON IDENT EOF | exp0 EOF | EOF
file : MODULE IDENT WHERE line* EOF
path : IDENT
line : IMPORT path+ | OPTION IDENT IDENT | declarations
ident : IRRREF | IDENT
lense : LPARENS ident+ COLON exp1 RPARENS
telescope : lense telescope
params : telescope | []
declarations :
  | DEF IDENT params DEFEQ exp1
  | DEF IDENT params COLON exp1 DEFEQ exp1
  | AXIOM IDENT params COLON exp1
face : LPARENS IDENT IDENT IDENT RPARENS
part : face+ ARROW exp1
exp0 : exp1 COMMA exp0 | exp1
exp1 : LSQ separated(COMMA, part) RSQ
  | LAM telescope COMMA exp1 | PI telescope COMMA exp1
  | SIGMA telescope COMMA exp1 | LSQ IRRREF ARROW exp1 RSQ
  | LT ident+ GT exp1 | exp2 ARROW exp1
  | exp2 PROD exp1 | exp2
exp2 : exp2 exp3 | exp2 APPFORMULA exp3 | exp3
exp3 : LPARENS exp0 RPARENS LSQ exp0 MAP exp0 RSQ
  | HOLE | PRE | KAN | IDJ exp3
  | exp3 FST | exp3 SND | NEGATE exp3 | INC exp3
  | exp3 AND exp3 | exp3 OR exp3 | ID exp3 | REF exp3
  | OUC exp3 | PATHP exp3 | PARTIAL exp3 | IDENT
  | IDENT LSQ exp0 MAP exp0 RSQ | HCOMP exp3
  | LPARENS exp0 RPARENS | TRANSP exp3 exp3

|<>\rightarrow module import data split where comp fill glue glue
unglue.1.2,
cubical

```

$\text{Per} \mathcal{U}_i \Pi \tau = \text{typePerAnders} \Sigma \tau \Pi$

$\mathbb{N} \omega$

$\mathbb{N} \text{Anders} \omega = \{V_i, U_i\}$

$= \alpha \beta$

$= V_i U_i \Pi$

$\tau$

$\tau = \{\text{infer, app, check, act, conv, eval}\}$

$\Sigma$

$\Sigma$

$\mathcal{P}$

$\int = \{\Pi, \Sigma, =, \mathbf{W}, \mathbf{0}, \mathbf{1}, \mathbf{2}, \mathbf{Path}, \mathbf{Glue}\}$

$e_i : O_{n+1} \rightarrow O_n O_{\text{CPS}} = O_0$   
 $O_{\text{PTS}}^9 O_{\text{CCHM}}^{10}$

---

---

$\Pi\Sigma=$

123

$4_{\infty}$

---



$\perp$   
 $\top$   
 $A \vee B$   
 $A \wedge B$   
 $A \Rightarrow B$   
 $x) \exists x:A B(x)$   
 $x) \forall x:A B(x)$   
 $=_A$        $A^I$

---

PerAndersAndersAnders

Anders

AndersPer

Anders

ΠΣΙδ

ΠΣ

$\Pi\Sigma$

$\Pi$

$\Pi\Pi$

$\Pi\Pi f : \Pi(x : A), B(x) \rightarrow U$

$$\Pi : U =_{\text{def}} \prod_{x:A} B(x).$$

`def Pi (A : U) (B : A → U) : U := Π (x : A), B(x)`

$\Pi\lambda x. b(x) \rightarrow b(x)$

$$\lambda(x : A) \rightarrow b(x) =_{\text{def}} \prod_{A:U} \prod_{B:A \rightarrow U} \prod_{b:\prod_{a:A} B(a)} \lambda x. b(x) : \prod_{y:A} B(y).$$

`def lambda (A : U) (B : A → U) (b : Π A B) : Π A B := λ (x : A), b x`

`def lam (A B : U) (f : A → B) : A → B := λ (x : A), f x`

$f : A \rightarrow B$

ΠΠΠΠλ

$$f\mathbf{a} =_{\text{def}} \prod_{A:\mathcal{U}} \prod_{B:A \rightarrow \mathcal{U}} \prod_{a:A} \prod_{f:\prod_{x:A} B(a)} f(\mathbf{a}) : B(\mathbf{a}).$$

```
def apply (A B: U) (f: A → B) (a: A) : B := f a
def app (A: U) (B: A → U) (a: A) (f: Pi A B): B a := f a
```

```
def oT (a b c: U) : U := (b → c) → (a → b) → (a → c)
def o (a b c : U) : oT a b c := λ (g: b → c) (f: a → b) (x: a), g (f x)
```

Πβlam ∘ app

$$f(\mathbf{a}) =_{B(\mathbf{a})} (\lambda(x:A) \rightarrow f(\mathbf{a}))(\mathbf{a}).$$

```
def Π-β (A : U) (B : A → U) (a : A) (f : Pi A B)
  : Path (B a) (Π-apply A B (Π-lambda A B f) a) (f a)
:= idp (B a) (f a)
```

Πηapp ∘ lam

$$f =_{(x:A) \rightarrow B(\mathbf{a})} (\lambda(y:A) \rightarrow f(\mathbf{y})).$$

```
def Π-η (A : U) (B : A → U) (a : A) (f : Pi A B)
  : Path (Pi A B) f (λ (x : A), f x)
:= idp (Pi A B) f
```

ΠΠ

$$p : E \rightarrow B \quad y : B \quad x : E \quad p(x) = y$$

$$F \rightarrow E \xrightarrow{p} BEFB(F, E, p, B) \quad p : E \rightarrow B \quad y : B \quad \text{B} \times F : p^{-1}(U_b) \rightarrow U_b \times F$$

$$p^{-1}(U_b) \times F \xrightarrow{\text{pr}_1} U_b \times F$$

$$FBapp : F \times B \rightarrow E$$

$$F \times B \xrightarrow{\text{app}} E \xrightarrow{\text{pr}_1} B$$

$$\text{pr}_1 \text{pr}_1 \text{app} \text{Set}_{/B} F A A \times B \rightarrow E \text{Set}_{/B} A \rightarrow F$$

$$E \Sigma(B, F) p = \text{pr}_1(F, \Sigma(B, F), \text{pr}_1, B)$$

$$(F, B * F, \text{pr}_1, B) y : BF(y)$$

FiberPi (B: U) (F: B -> U) (y: B)

: Path U (fiber (Sigma B F) B (pi1 B F) y) (F y)

ΠΣ

$$g : B \rightarrow A \quad C \Pi_g : C_{/B} \rightarrow C_{/A}$$

$$H(\infty, 1) E \rightarrow B : H_{/B} H \Gamma_\Sigma(E)$$

$$\Gamma_\Sigma(E) = \Pi_\Sigma(E) \in H.$$

setFun (A B : U) ( \_: isSet B) : isSet (A -> B)

piIsContr (A: U) (B: A -> U) (u: isContr A)

(q: (x: A) -> isContr (B x)) : isContr (Pi A B)

$$f : A \rightarrow B \quad g : B \rightarrow A \quad f \circ g : B \xrightarrow{g} A \xrightarrow{f} B$$

$\Sigma$

$\Sigma\Sigma$

$\Sigma\Sigma AB\cup A : \cup B : A \rightarrow \cup a : AB(a)\Sigma(A, B)(x, B(x))$

```
def Sigma (A : U) (B : A -> U) : U :=  $\Sigma$  (x : A), B x
```

$\Sigma(a, b)\Sigma(A, B)\Sigma(A, B)(a, b)b : B(a)$

```
def dpair (A : U) (B : A -> U) (a : A) (b : B a) : Sigma A B = (a,b)
```

$\Sigma\Sigma$

```
def pr1 (A : U) (B : A -> U) (x : Sigma A B) : A := x.1
def pr2 (A : U) (B : A -> U) (x : Sigma A B) : B (pr1 A B x) := x.2
def sigInd (A : U) (B : A -> U) (C : Sigma A B -> U)
  (g : (a : A) (b : B a) -> C (a, b)) (p : Sigma A B)
  : C p := g p.1 p.2
```

$\Sigma pr_1 pr_2(a, b)abEqu$

```
def Beta1 (A : U) (B : A -> U) (a : A) (b : B a)
  : Equ A a (pr1 A B (a,b))
```

```
def Beta2 (A : U) (B : A -> U) (a : A) (b : B a)
  : Equ (B a) b (pr2 A B (a,b))
```

$\Sigma p : \Sigma(A, B)(pr_1, pr_2)$

```
def Eta2 (A : U) (B : A -> U) (p : Sigma A B)
  : Equ (Sigma A B) p (pr1 A B p, pr2 A B p)
```

$f : A \rightarrow BC\Sigma_f : C/A \rightarrow C/B$

$x : Ay : BR(x, y)f : A \rightarrow Bx : AR(x, f(x))$

```
ac (A B : U) (R : A -> B -> U)
  : (p : (x:A) -> (y:B)*(R x y)) -> (f:A->B) * ((x:A)->R(x)(f x))
```

```
total (A:U) (B C : A -> U)
  (f : (x:A) -> B x -> C x) (w : Sigma A B)
  : Sigma A C = (w.1, f (w.1) (w.2))
```

$\Sigma\Sigma$

```
setSig (A:U) (B : A -> U) (sA : isSet A)
  (sB : (x:A) -> isSet (B x)) : isSet (Sigma A B)
```

$t, u : \Sigma(A, B) \rightarrow t_1 =_A u_1 \rightarrow t_2 =_{B(p@i)} u_2$

```
pathSig (A:U) (B : A -> U) (t u : Sigma A B)
  : Path U (Path (Sigma A B) t u)
    ((p: Path A t.1 u.1) * PathP (<i>B(p@i)) t.2 u.2)
```

[0, 1]<sup>56789</sup>

Hetero (A B: U) (a: A) (b: B) (P: Path U A B) : U = PathP P a b  
Path (A: U) (a b: A) : U = PathP (<i>A) a b

[0, 1]<i>aλ(i : I) → a

refl (A: U) (a: A) : Path A a a

app1 (A: U) (a b: A) (p: Path A a b): A = p @ 0

app2 (A: U) (a b: A) (p: Path A a b): A = p @ 1

λ(i : I) → ~~comp~~  
~~Path~~  
~~A~~  
~~a~~  
~~b~~

composition (A: U) (a b c: A) (p: Path A a b) (q: Path A b c)  
: Path A a c = comp (<i>Path A a (q@i)) p []

inv (A: U) (a b: A) (p: Path A a b): Path A b a = <i>p @ -i

λ(i, j : I) → p@min(i, j)λ(i, j : I) → p@max(i, j)

λ(i : I)λ(i : I) → a ~~Path~~  
~~A~~  
~~a~~  
~~b~~ λ(i : I) → b  
~~Path~~  
~~A~~  
~~a~~  
~~b~~

connection1 (A: U) (a b: A) (p: Path A a b)  
: PathP (<x> Path A (p@x) b) p (<i>b)  
= <y x> p @ (x \ / y)

connection2 (A: U) (a b: A) (p: Path A a b)  
: PathP (<x> Path A a (p@x)) (<i>a) p  
= <x y> p @ (x / \ y)

[0, 1]λ→

- 
- <https://5ht.co/cubicaltt.pdf>
  - <https://5ht.co/ccctt.pdf>
  - <http://www.cse.chalmers.se/~coquand/mod1.pdf>

- <https://www.cs.cmu.edu/~cangiuli/papers/ccctt.pdf>
- <https://arxiv.org/pdf/1712.04864.pdf>

```
ap (A B: U) (f: A -> B)
  (a b: A) (p: Path A a b)
  : Path B (f a) (f b)
```

```
apd (A: U) (a x:A) (B: A -> U) (f: A -> B a)
  (b: B a) (p: Path A a x)
  : Path (B a) (f a) (f x)
```

p

```
trans (A B: U) (p: Path U A B) (a: A) : B
```

```
singl (A: U) (a: A): U = (x: A) * Path A a x
```

```
eta (A: U) (a: A): singl A a = (a, refl A a)
```

```
contr (A: U) (a b: A) (p: Path A a b)
  : Path (singl A a) (eta A a) (b,p)
  = <i> (p @ i, <j> p @ i/\j)
```

```
D (A: U) : U = (x y: A) -> Path A x y -> U
```

```
J (A: U) (x y: A) (C: D A)
  (d: C x x (refl A x))
  (p: Path A x y) : C x y p
= subst (singl A x) T (eta A x) (y, p) (contr A x y p) d where
  T (z: singl A x) : U = C x (z.1) (z.2)
```

```
J (A: U) (a b: A)
  (P: singl A a -> U)
  (u: P (a, refl A a))
  (p: Path A a b) : P (b,p)
```

```
J (A: U) (a b: A)
  (C: (x: A) -> Path A a x -> U)
  (d: C a (refl A a))
  (p: Path A a b) : C b p
```

```
trans_comp (A: U) (a: A)
  : Path A a (trans A A (<_> A) a)
  = fill (<i> A) a []
```

```
subst_comp (A: U) (P: A -> U) (a: A) (e: P a)
  : Path (P a) e (subst A P a a (refl A a) e)
  = trans_comp (P a) e
```

```
J_comp (A: U) (a: A) (C: (x: A) -> Path A a x -> U) (d: C a (refl A a))
  : Path (C a (refl A a)) d (J A a C d a (refl A a))
  = subst_comp (singl A a) T (eta A a) d where T (z: singl A a)
  : U = C a (z.1) (z.2)
```

$$\mathcal{U}_{n \in \mathbb{N}} \mathcal{U}_0$$
$$\mathcal{U}_i : \mathcal{U}_j, i, j \in \mathbb{N}, j$$
$$\mathcal{U}_i \rightarrow \mathcal{U}_j : \mathcal{U}_{\lambda(i,j), i, j \in \mathbb{N}} \lambda : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$$
$$\lambda_{\max} : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$$
$$\lambda_{\text{snd}} : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$$
$$\mathcal{U}_i, i \in \mathbb{N} \in$$
$$\mathcal{U}_{n \in \mathbb{N}} \mathcal{U}_i : \mathcal{U}_j, i, j \in \mathbb{N} \mathcal{U}_i \rightarrow \mathcal{U}_j : \mathcal{U}_{\lambda(i,j), i, j \in \mathbb{N}} \lambda_{\max} \lambda_{\text{snd}}$$
$$\{\{\star, \square\}, \{\star : \square\}, \{i \rightarrow j : j; i, j \in \{\star, \square\}\} \star \square \lambda = \text{snd}$$
$$\text{PTS}^\infty \{\mathcal{U}_{i \in \mathbb{N}}, \mathcal{U}_i : \mathcal{U}_j; i < j; i, j \in \mathbb{N}, \mathcal{U}_i \rightarrow \mathcal{U}_j : \mathcal{U}_{\lambda(i,j); i, j \in \mathbb{N}}\} \mathcal{U}_i \text{iiPTS}^{\infty 10}$$
$$\gamma_0 : \Gamma =_{\text{def}} \star.$$
$$\Gamma; A =_{\text{def}} \sum_{\gamma: \Gamma} A(\gamma).$$
$$\Gamma \vdash A =_{\text{def}} \prod_{\gamma: \Gamma} A(\gamma).$$

$\Pi\Sigma$

$\Pi\Sigma\Sigma$

```
def MLTT (A: U) : U :=  $\Sigma$ 
  ( $\Pi$ -form :  $\Pi$  (B: A  $\rightarrow$  U), U)
  ( $\Pi$ -ctor1 :  $\Pi$  (B: A  $\rightarrow$  U),  $\Pi$  A B  $\rightarrow$   $\Pi$  A B)
  ( $\Pi$ -elim1 :  $\Pi$  (B: A  $\rightarrow$  U),  $\Pi$  A B  $\rightarrow$   $\Pi$  A B)
  ( $\Pi$ -comp1 :  $\Pi$  (B: A  $\rightarrow$  U) (a: A) (f:  $\Pi$  A B),
    Equ (B a) ( $\Pi$ -elim1 B ( $\Pi$ -ctor1 B f) a) (f a))
  ( $\Pi$ -comp2 :  $\Pi$  (B: A  $\rightarrow$  U) (a: A) (f:  $\Pi$  A B),
    Equ ( $\Pi$  A B) f ( $\lambda$  (x : A), f x))
  ( $\Sigma$ -form :  $\Pi$  (B: A  $\rightarrow$  U), U)
  ( $\Sigma$ -ctor1 :  $\Pi$  (B: A  $\rightarrow$  U) (a : A) (b : B a),  $\Sigma$  A B)
  ( $\Sigma$ -elim1 :  $\Pi$  (B: A  $\rightarrow$  U) (p :  $\Sigma$  A B), A)
  ( $\Sigma$ -elim2 :  $\Pi$  (B: A  $\rightarrow$  U) (p :  $\Sigma$  A B), B (pr1 A B p))
  ( $\Sigma$ -comp1 :  $\Pi$  (B: A  $\rightarrow$  U) (a : A) (b: B a),
    Equ A a ( $\Sigma$ -elim1 B ( $\Sigma$ -ctor1 B a b)))
  ( $\Sigma$ -comp2 :  $\Pi$  (B: A  $\rightarrow$  U) (a : A) (b: B a),
    Equ (B a) b ( $\Sigma$ -elim2 B (a, b)))
  ( $\Sigma$ -comp3 :  $\Pi$  (B: A  $\rightarrow$  U) (p :  $\Sigma$  A B),
    Equ ( $\Sigma$  A B) p (pr1 A B p, pr2 A B p))
  ( $=$ -form :  $\Pi$  (a: A), A  $\rightarrow$  U)
  ( $=$ -ctor1 :  $\Pi$  (a: A), Equ A a a)
  ( $=$ -elim1 :  $\Pi$  (a: A) (C: D A) (d: C a a ( $=$ -ctor1 a))
    (y: A) (p: Equ A a y), C a y p)
  ( $=$ -comp1 :  $\Pi$  (a: A) (C: D A) (d: C a a ( $=$ -ctor1 a)),
    Equ (C a a ( $=$ -ctor1 a)) d ( $=$ -elim1 a C d a ( $=$ -ctor1 a))), U
```

```
theorem instance (A : U) : MLTT A :=
  ( $\Pi$  A, lambda A, app A, comp1 A, comp2 A,
   $\Sigma$  A, pair A, pr1 A, pr2 A, comp3 A, comp4 A, comp5 A,
  Equ A, refl A, J A, comp6 A, A)
```

mltt.cttcubicaltt

\$ rlwrap ./anders.native check ./experiments/mltt.anders

File loaded.

> :n instance

TYPE:  $\prod (A : U), \Sigma (\prod\text{-form} : \prod (B : (A \rightarrow U)), U), \Sigma (\prod\text{-ctor}_1 : \prod (B : (A \rightarrow U)), (\prod (x : A), (B x) \rightarrow \prod (x : A), (B x))), \Sigma (\prod\text{-elim}_1 : \prod (B : (A \rightarrow U)), (\prod (x : A), (B x) \rightarrow \prod (x : A), (B x))), \Sigma (\prod\text{-comp}_1 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (f : \prod (x : A), (B x)), \prod (P : ((B a) \rightarrow U)), ((P ((\prod\text{-elim}_1 B) ((\prod\text{-ctor}_1 B) f)) a)) \rightarrow (P (f a))), \Sigma (\prod\text{-comp}_2 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (f : \prod (x : A), (B x)), \prod (P : (\prod (x : A), (B x) \rightarrow U)), ((P f) \rightarrow (P \lambda (x : A), (f x)))), \Sigma (\Sigma\text{-form} : \prod (B : (A \rightarrow U)), U), \Sigma (\Sigma\text{-ctor}_1 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (b : (B a)), \Sigma (x : A), (B x)), \Sigma (\Sigma\text{-elim}_1 : \prod (B : (A \rightarrow U)), \prod (p : \Sigma (x : A), (B x)), A), \Sigma (\Sigma\text{-elim}_2 : \prod (B : (A \rightarrow U)), \prod (p : \Sigma (x : A), (B x)), (B p .1)), \Sigma (\Sigma\text{-comp}_1 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (b : (B a)), \prod (P : (A \rightarrow U)), ((P a) \rightarrow (P ((\Sigma\text{-elim}_1 B) ((\Sigma\text{-ctor}_1 B) a) b)))), \Sigma (\Sigma\text{-comp}_2 : \prod (B : (A \rightarrow U)), \prod (a : A), \prod (b : (B a)), \prod (P : ((B a) \rightarrow U)), ((P b) \rightarrow (P ((\Sigma\text{-elim}_2 B) (a, b))))), \Sigma (\Sigma\text{-comp}_3 : \prod (B : (A \rightarrow U)), \prod (p : \Sigma (x : A), (B x)), \prod (P : (\Sigma (x : A), (B x) \rightarrow U)), ((P p) \rightarrow (P (p .1, p.2)))), \Sigma (=form : \prod (a : A), (A \rightarrow U)), \Sigma (=ctor_1 : \prod (a : A), \prod (P : (A \rightarrow U)), ((P a) \rightarrow (P a))), \Sigma (=elim_1 : \prod (a : A), \prod (C : \prod (x : A), \prod (y : A), (\prod (P : (A \rightarrow U)), ((P x) \rightarrow (P y)) \rightarrow U)), \prod (d : ((C a) a) (=ctor_1 a))), \prod (y : A), \prod (p : \prod (P : (A \rightarrow U)), ((P a) \rightarrow (P y))), ((C a) y) p)), \Sigma (=comp_1 : \prod (a : A), \prod (C : \prod (x : A), \prod (y : A), (\prod (P : (A \rightarrow U)), ((P x) \rightarrow (P y)) \rightarrow U)), \prod (d : (((C a) a) (=ctor_1 a))), \prod (P : (((C a) a) (=ctor_1 a)) \rightarrow U)), ((P d) \rightarrow (P (((=elim_1 a) C) d) a) (=ctor_1 a))))), U$

NORMEVAL:  $\lambda (A : U), (\lambda (B : (A \rightarrow U)), \prod (x : A), (B x), (\lambda (B : (A \rightarrow U)), \lambda (b : \prod (x : A), (B x)), \lambda (x : A), (b x), (\lambda (B : (A \rightarrow U)), \lambda (f : \prod (x : A), (B x)), \lambda (a : A), (f a), (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (f : \prod (x : A), (B x)), \lambda (P : ((B a) \rightarrow U)), \lambda (u : (P (f a))), u, (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (f : \prod (x : A), (B x)), \lambda (P : (\prod (x : A), (B x) \rightarrow U)), \lambda (u : (P f)), u, (\lambda (B : (A \rightarrow U)), \Sigma (x : A), (B x), (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (b : (B a)), (a, b), (\lambda (B : (A \rightarrow U)), \lambda (x : \Sigma (x : A), (B x)), x.1, (\lambda (B : (A \rightarrow U)), \lambda (x : \Sigma (x : A), (B x)), x.2, (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (b : (B a)), \lambda (P : (A \rightarrow U)), \lambda (u : (P a)), u, (\lambda (B : (A \rightarrow U)), \lambda (a : A), \lambda (b : (B a)), \lambda (P : ((B a) \rightarrow U)), \lambda (u : (P b)), u, (\lambda (B : (A \rightarrow U)), \lambda (p : \Sigma (x : A), (B x)), \lambda (P : (\Sigma (x : A), (B x) \rightarrow U)), \lambda (u : (P p)), u, (\lambda (x : A), \lambda (y : A), \prod (P : (A \rightarrow U)), ((P x) \rightarrow (P y))), (\lambda (x : A), \lambda (P : (A \rightarrow U)), \lambda (u : (P x)), u, ((J A), ((comp_6 A), A))))))))))))))$

```
data empty =
emptyRec (C: U): empty -> C = split {}
emptyInd (C: empty -> U): (z: empty) -> C z = split {}

data unit = star
unitRec (C: U) (x: C): unit -> C = split tt -> x
unitInd (C: unit -> U) (x: C tt): (z: unit) -> C z = split tt -> x
```

```

data bool = false | true
b1: U = bool -> bool
b2: U = bool -> bool -> bool
negation: b1 = split { false -> true; true -> false }
or: b2 = split { false -> idfun bool; true -> lambda bool bool true }
and: b2 = split { false -> lambda bool bool false; true -> idfun bool }
boolEq: b2 = lambda bool (bool -> bool) negation
boolRec (C: U) (f t: C): bool -> C = split { false -> f ; true -> t }
boolInd (C: bool -> U) (f: A false) (t: A true): (n:bool) -> A n
    = split { false -> f ; true -> t }

```

$$M_A(X) = 1 + A$$

```

data maybe (A: U) = nothing | just (x: A)
maybeRec (A P: U) (n: P) (j: A -> P): maybe A -> P
    = split { nothing -> n; just a -> j a }

maybeInd (A: U) (P: maybe A -> U) (n: P nothing)
    (j: (a: A) -> P (just a)): (a: maybe A) -> P a
    = split { nothing -> n ; just x -> j x }

```

```

data either (A B: U) = left (x: A) | right (y: B)
eitherRec (A B C: U) (b: A -> C) (c: B -> C): either A B -> C
    = split { inl x -> b(x) ; inr y -> c(y) }

```

```

eitherInd (A B: U) (C: either A B -> U)
    (x: (a: A) -> C (inl a))
    (y: (b: B) -> C (inr b))
: (x: either A B) -> C x
    = split { inl i -> x i ; inr j -> y j }

```

```

data tuple (A B: U) = pair (x: A) (y: B)
prod (A B: U) (x: A) (y: B): (_, A) * B = (x,y)
tupleRec (A B C: U) (c: (x:A) (y:B) -> C): (x: tuple A B) -> C
    = split pair a b -> c a b
tupleInd (A B: U) (C: tuple A B -> U)
    (c: (x:A)(y:B) -> C (pair x y))
: (x: tuple A B) -> C x
    = split pair a b -> c a b

```

```

data nat = zero | succ (n: nat)
natEq: nat -> nat -> bool
natCase (C:U) (a b: C): nat -> C
natRec (C:U) (z: C) (s: nat->C->C) : (n:nat) -> C

```

```

natElim (C:nat->U) (z: C zero)
  (s: (n:nat)->C(succ n)): (n:nat) -> C(n)
natInd (C:nat->U) (z: C zero)
  (s: (n:nat)->C(n)->C(succ n)): (n:nat) -> C(n)

```

$(\mu L_A, \text{in})L_A(X) = 1 + (A \times X)\mu L_A = \text{List}(A)\text{nil} : 1 \rightarrow \text{List}(A)\text{cons} : A \times \text{List}(A) \rightarrow \text{List}(A)\text{nil} = \text{in} \circ \text{inlcons} = \text{in} \circ \text{inr} \text{in} = [\text{nil}, \text{cons}]$

```

data list (A: U) = nil | cons (x:A) (xs: list A)
listCase (A C:U) (a b: C): list A -> C
listRec (A C:U) (z: C) (s: A->list A->C->C): (n:list A) -> C
listElim (A: U) (C:list A->U) (z: C nil)
  (s: (x:A)(xs:list A)->C(cons x xs)): (n:list A) -> C(n)
listInd (A: U) (C:list A->U) (z: C nil)
  (s: (x:A)(xs:list A)->C(xs)->C(cons x xs)): (n:list A) -> C(n)

```

```

null (A:U): list A -> bool
head (A:U): list A -> maybe A
tail (A:U): list A -> maybe (list A)
nth (A:U): nat -> list A -> maybeA
append (A: U): list A -> list A -> list A
reverse (A: U): list A -> list A
map (A B: U): (A -> B) -> list A -> list B
zip (AB: U): list A -> list B -> list (tuple A B)
foldr (AB: U): (A -> B -> B) -> B -> list A -> B
foldl (AB: U): (B -> A -> B) -> B -> list A -> B
switch (A: U): (Unit -> list A) -> bool -> list A
filter (A: U): (A -> bool) -> list A -> list A
length (A: U): list A -> nat
listEq (A: eq): list A.1 -> list A.1 -> bool

```

```

data stream (A: U) = cons (x: A) (xs: stream A)

```

```

data fin (n: nat)
  = fzero | fsucc (_, fin (pred n))

```

```

fz (n: nat): fin (succ n) = fzero
fs (n: nat): fin n -> fin (succ n) = \ (x: fin n) -> fsucc x

```

```

data vector (A: U) (n: nat)
  = nil | cons (_: A) (_: vector A (pred n))

data seq (A: U) (B: A -> A -> U) (X Y: A)
  = seqNil (_: A)
  | seqCons (X Y Z: A) (_: B X Y) (_: Seq A B Y Z)

nat = (X:U) -> (X -> X) -> X -> X
      (X -> X) succ X zero

list (A: U) = (X:U) -> X -> (A -> X) -> X

NAT (A: U) = (X:U) -> isSet X -> X -> (A -> X) -> X

TRUN (A:U) type = (X: U) -> isProp X -> (A -> X) -> X
S1 = (X:U) -> isGroupoid X -> ((x:X) -> Path X x x) -> X
MONOPL (A:U) = (X:U) -> isSet X -> (A -> X) -> X
NAT = (X:U) -> isSet X -> X -> (A -> X) -> X

upPath (X Y:U) (f:X->Y) (a:X->X): X -> Y = o X X Y f a
downPath (X Y:U) (f:X->Y) (b:Y->Y): X -> Y = o X Y Y b f
naturality (X Y:U) (f:X->Y) (a:X->X) (b:Y->Y): U
  = Path (X->Y) (upPath X Y f a) (downPath X Y f b)

unitEnc': U = (X: U) -> isSet X -> X -> X
isUnitEnc (one: unitEnc'): U
  = (X Y:U) (x:isSet X) (y:isSet Y) (f:X->Y) ->
    naturality X Y f (one X x) (one Y y)

unitEnc: U = (x: unitEnc') * isUnitEnc x
unitEncStar: unitEnc = (\(X:U) (_:isSet X) ->
  idfun X, \ (X Y: U) (_:isSet X) (_:isSet Y) -> refl (X->Y))
unitEncRec (C: U) (s: isSet C) (c: C): unitEnc -> C
  = \ (z: unitEnc) -> z.1 C s c
unitEncBeta (C: U) (s: isSet C) (c: C)
  : Path C (unitEncRec C s c unitEncStar) c = refl C c
unitEncEta (z: unitEnc): Path unitEnc unitEncStar z = undefined
unitEncInd (P: unitEnc -> U) (a: unitEnc): P unitEncStar -> P a
  = subst unitEnc P unitEncStar a (unitEncEta a)
unitEncCondition (n: unitEnc'): isProp (isUnitEnc n)
  = \ (f g: isUnitEnc n) ->
    <h> \ (x y: U) -> \ (X: isSet x) -> \ (Y: isSet y)
    -> \ (F: x -> y) -> <i> \ (R: x) -> Y (F (n x X R)) (n y Y (F R))
    (<j> f x y X Y F @ j R) (<j> g x y X Y F @ j R) @ h @ i

```

∞

$\mathbb{R}^{n11}$

$I = [0, 1]$

```
data I = i0
      | i1
      | seg <i> [(i=0) -> i0,
                (i=1) -> i1]
```

$i_0, i_1 : x, y : A$

I

```
pathToHtpy (A: U) (x y: A) (p: Path A x y): I -> A
= split { i0 -> x; i1 -> y; seg @ i -> p @ i }
```

$f, g : X \rightarrow Y, H : X \times I \rightarrow Y$

$$\begin{cases} H(x, 0) = f(x), \\ H(x, 1) = g(x). \end{cases}$$

```
homotopy (X Y: U) (f g: X -> Y)
(p: (x: X) -> Path Y (f x) (g x))
(x: X): I -> Y = pathToHtpy Y (f x) (g x) (p x)
```

1213

```
cat: U = (A: U) * (A -> A -> U)
groupoid: U = (X: cat) * isCatGroupoid X
PathCat (X: U): cat = (X, \ (x y: X) -> Path X x y)
```

---

R

<http://www.cse.chalmers.se/~coquand/Proposal.pdf>

```

isCatGroupoid (C: cat): U
= (id: (x: C.1) -> C.2 x x)
* (c: (x y z:C.1) -> C.2 x y -> C.2 y z -> C.2 x z)
* (inv: (x y: C.1) -> C.2 x y -> C.2 y x)
* (inv_left: (x y: C.1) (p: C.2 x y) ->
  Path (C.2 x x) (c x y x p (inv x y p)) (id x))
* (inv_right: (x y: C.1) (p: C.2 x y) ->
  Path (C.2 y y) (c y x y (inv x y p) p) (id y))
* (left: (x y: C.1) (f: C.2 x y) ->
  Path (C.2 x y) (c x x y (id x) f) f)
* (right: (x y: C.1) (f: C.2 x y) ->
  Path (C.2 x y) (c x y y f (id y)) f)
* ((x y z w:C.1)(f:C.2 x y)(g:C.2 y z)(h:C.2 z w)->
  Path (C.2 x w) (c x z w (c x y z f g) h)
    (c x y w f (c y z w g h)))

```

```

PathGrpd (X: U)
: groupoid
= ((Ob,Hom),id,c,sym X,compPathInv X,compInvPath X,L,R,Q) where
Ob: U = X
Hom (A B: Ob): U = Path X A B
id (A: Ob): Path X A A = refl X A
c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
= comp <i> Path X A (g@i) f []

```

$\Pi\Sigma$

```

funext_form (A B: U) (f g: A -> B): U
= Path (A -> B) f g

```

```

funext (A B: U) (f g: A -> B) (p: (x:A) -> Path B (f x) (g x))
: funext_form A B f g
= <i> \ (a: A) -> p a @ i

```

```

happly (A B: U) (f g: A -> B) (p: funext_form A B f g) (x: A)
: Path B (f x) (g x)
= cong (A -> B) B (\ (h: A -> B) -> apply A B h x) f g p

```

```

funext_Beta (A B: U) (f g: A -> B) (p: (x:A) -> Path B (f x) (g x))
: (x:A) -> Path B (f x) (g x)
= \ (x:A) -> haply A B f g (funext A B f g p) x

```

```

funext_Eta (A B: U) (f g: A -> B) (p: Path (A -> B) f g)
: Path (Path (A -> B) f g) (funext A B f g (happly A B f g p)) p
= refl (Path (A -> B) f g) p

```

```

pullback (A B C:U) (f: A -> C) (g: B -> C): U
= (a: A)
* (b: B)
* Path C (f a) (g b)

pb1 (A B C: U) (f: A -> C) (g: B -> C)
: pullback A B C f g -> A
= \x: pullback A B C f g -> x.1

pb2 (A B C: U) (f: A -> C) (g: B -> C)
: pullback A B C f g -> B
= \x: pullback A B C f g -> x.2.1

pb3 (A B C: U) (f: A -> C) (g: B -> C)
: (x: pullback A B C f g) -> Path C (f x.1) (g x.2.1)
= \x: pullback A B C f g -> x.2.2

kernel (A B: U) (f: A -> B): U
= pullback A A B f f

hofiber (A B: U) (f: A -> B) (y: B): U
= pullback A unit B f (\x: unit) -> y

pullbackSq (Z A B C: U) (f: A -> C) (g: B -> C) (z1: Z -> A) (z2: Z -> B): U
= (h: (z:Z) -> Path C ((o Z A C f z1) z) (((o Z B C g z2)) z))
* isEquiv Z (pullback A B C f g) (induced Z A B C f g z1 z2 h)

completePullback (A B C: U) (f: A -> C) (g: B -> C)
: pullbackSq (pullback A B C f g) A B C f g (pb1 A B C f g) (pb2 A B C f
g)

data pushout (A B C: U) (f: C -> A) (g: C -> B)
= po1 (_: A)
| po2 (_: B)
| po3 (c: C) <i> [ (i = 0) -> po1 (f c) ,
(i = 1) -> po2 (g c) ]

```

```

isFBundle1 (B: U) (p: B → U) (F: U): U
= (λ (b: B) → isContr (Path U (p b) F))
* ((x: Sigma B p) → B)

isFBundle2 (B: U) (p: B → U) (F: U): U
= (V: U)
* (v: surjective V B)
* ((x: V) → Path U (p (v.1 x)) F)

im1 (A B: U) (f: A → B): U = (b: B) * pTrunc ((a:A) * Path B (f a) b)
BAut (F: U): U = im1 unit U (λ(x: unit) → F)
unitIm1 (A B: U) (f: A → B): im1 A B f → B = λ(x: im1 A B f) → x.1
unitBAut (F: U): BAut F → U = unitIm1 unit U (λ(x: unit) → F)

isFBundle3 (E B: U) (p: E → B) (F: U): U
= (X: B → BAut F)
* (classify B (BAut F) (λ(b: B) → fiber E B p b) (unitBAut F) X) where
classify (A' A: U) (E': A' → U) (E: A → U) (f: A' → A): U
= (x: A') → Path U (E'(x)) (E(f(x)))

isFBundle4 (E B: U) (p: E → B) (F: U): U
= (V: U)
* (v: surjective V B)
* (v': prod V F → E)
* pullbackSq (prod V F) E V B p v.1 v' (λ(x: prod V F) → x.1)

fiber (A B: U) (f: A → B) (y: B): U = (x: A) * Path B y (f x)
isSingleton (X:U): U = (c:X) * ((x:X) → Path X c x)
isEquiv (A B: U) (f: A → B): U = (y: B) → isContr (fiber A B f y)
equiv (A B: U): U = (f: A → B) * isEquiv A B f

isSurjective (A B: U) (f: A → B): U
= (b: B) * pTrunc (fiber A B f b)

surjective (A B: U): U
= (f: A → B)
* isSurjective A B f

isInjective' (A B: U) (f: A → B): U
= (b: B) → isProp (fiber A B f b)

injective (A B: U): U
= (f: A → B)
* isInjective A B f

```

```
isEmbedding (A B: U) (f: A -> B) : U
= (x y: A) -> isEquiv (Path A x y) (Path B (f x) (f y)) (cong A B f x y)
```

```
embedding (A B: U): U
= (f: A -> B)
* isEmbedding A B f
```

```
isHae (A B: U) (f: A -> B): U
= (g: B -> A)
* (eta_: Path (id A) (o A B A g f) (idfun A))
* (eps_: Path (id B) (o B A B f g) (idfun B))
* ((x: A) -> Path B (f ((eta_ @ 0) x)) ((eps_ @ 0) (f x)))
```

```
hae (A B: U): U
= (f: A -> B)
* isHae A B f
```

```
iso_Form (A B: U): U = isIso A B -> Path U A B
```

```
iso_Intro (A B: U): iso_Form A B
```

```
iso_Elim (A B: U): Path U A B -> isIso A B
```

```
iso_Comp (A B : U) (p : Path U A B)
: Path (Path U A B) (iso_Intro A B (iso_Elim A B p)) p
```

```
iso_Uniq (A B : U) (p: isIso A B)
: Path (isIso A B) (iso_Elim A B (iso_Intro A B p)) p
```

```
univ_Formation (A B: U): U = equiv A B -> Path U A B
```

```
equivToPath (A B: U): univ_Formation A B
= \ (p: equiv A B) -> <i> Glue B [(i=0) -> (A,p),
(i=1) -> (B, subst U (equiv B) B B (<_>B) (idEquiv B))] ]
```

```

pathToEquiv (A B : U) (p : Path U A B) : equiv A B
  = subst U (equiv A) A B p (idEquiv A)

```

```

eqToEq (A B : U) (p : Path U A B)
  : Path (Path U A B) (equivToPath A B (pathToEquiv A B p)) p
  = <j i> let Ai : U = p@i in Glue B
    [ (i=0) -> (A,pathToEquiv A B p),
      (i=1) -> (B,pathToEquiv B B (<k> B)),
      (j=1) -> (p@i,pathToEquiv Ai B (<k> p @ (i \ / k))) ]

```

```

transPathFun (A B : U) (w : equiv A B)
  : Path (A -> B) w.1 (pathToEquiv A B (equivToPath A B w)).1

```

```

data I = i0
  | i1
  | seg <i> [(i=0) -> i0,
            (i=1) -> i1]

```

$i_0, i_1 : x, y : A$

```

data S1
  = base
  | loop <i> [ (i = 0) -> base,
              (i = 1) -> base ]

```

```

data S2
  = point
  | surf <i j> [ (i = 0) -> point, (i = 1) -> point,
                (j = 0) -> point, (j = 1) -> point ]
                (j = 0) -> point, (j = 1) -> point ]

```

```

data susp (A: U)
  = north
  | south
  | merid (a: A) <i> [ (i = 0) -> north ,
                    (i = 1) -> south ]

```

```

data pTrunc (A: U) -- (-1)-trunc, mere proposition truncation
  = pinc (a: A)
  | pline (x y: pTrunc A) <i>
    [ (i = 0) -> x,
      (i = 1) -> y ]

```

```

data sTrunc (A: U) -- (0)-trunc, set truncation
  = sinc (a: A)
  | sline (a b: sTrunc A)
    (p q: Path (sTrunc A) a b) &lt;i>j&gt;;
    [ (i = 0) -> p @ j,
      (i = 1) -> q @ j,
      (j = 0) -> a,
      (j = 1) -> b ]

```

```

data gTrunc (A: U) -- (1)-trunc, groupoid truncation
  = ginc (a: A)
  | gline (a b: gTrunc A)
    (p q: Path (gTrunc A) a b)
    (r s: Path (Path (gTrunc A) a b) p q) &lt;i>j k&gt;;
    [ (i = 0) -> r @ j @ k,
      (i = 1) -> s @ j @ k,
      (j = 0) -> p @ k,
      (j = 1) -> q @ k,
      (k = 0) -> a,
      (k = 1) -> b ]

```

```

data quot (A: U) (R: A -> A -> U)
  = inj (a: A)
  | quoteq (a b: A) (r: R a b) &lt;i>&gt;;
    [ (i = 0) -> inj a,
      (i = 1) -> inj b ]

```

```

data setquot (A: U) (R: A -> A -> U)
  = quotient (a: A)
  | identification (a b: A) (r: R a b) &lt;i>&gt;;
    [ (i = 0) -> quotient a,
      (i = 1) -> quotient b ]

```

```

| setTruncation (a b: setquot A R)
  (p q: Path (setquot A R) a b) &lt;i j&gt;;
  [ (i = 0) -> p @ j,
    (i = 1) -> q @ j,
    (j = 0) -> a,
    (j = 1) -> b ]

```

```
storage: U -> U = list
```

$\Sigma$

```

process : U
= (protocol state: U)
* (current: prod protocol state)
* (act: id (prod protocol state))
* (storage (prod protocol state))

```

```

spawn (protocol state: U) (init: prod protocol state)
  (action: id (prod protocol state)) : process
= (protocol,state,init,action,nil)

```

```

protocol (p: process): U = p.1
state    (p: process): U = p.2.1
signature (p: process): U = prod p.1 p.2.1
current  (p: process):      signature p = p.2.2.1
action   (p: process):      id (signature p) = p.2.2.2.1
trace    (p: process):      storage (signature p) = p.2.2.2.2

```

$$P \times S \rightarrow SP \times S \rightarrow P \times S$$

```
receive (p: process) : protocol p = axiom
```

```
send (p: process) (message: protocol p) : unit = axiom
```

```

execute (p: process) (message: protocol p) : process
= let step: signature p = (action p) (message, (current p).2)
  in (protocol p, state p, step, action p, cons step (trace p))

```

14

15

16

---

$\Sigma_{A:U} A \rightarrow A \rightarrow \text{U} \text{Upr}_1 \text{Obpr}_2 \text{Hom}(a, b) a, b : \text{Ob}$

$\text{cat} : U = (A : U) * (A \rightarrow A \rightarrow U)$

$\text{CHom}_C(a, b) a, b : \text{Ob}_C \text{idHom}_C(x, x)$

$\text{COb}_C a, b : \text{Ob}_C \text{Hom}_C(a, b) a : \text{Ob}_C 1_a : \text{Hom}_C(a, a) a, b, c :$   
 $\text{Ob}_C \text{Hom}_C(b, c) \rightarrow \text{Hom}_C(a, b) \rightarrow \text{Hom}_C(a, c) g \circ f a, b : \text{Ob}_C f :$   
 $\text{Hom}_C(a, b) f = 1_b \circ f f = f \circ 1_a a, b, c, d : \text{Af} : \text{Hom}_C(a, b) g : \text{Hom}_C(b, c)$   
 $h : \text{Hom}_C(c, d) h \circ (g \circ f) = (h \circ g) \circ f$

$a, b : \text{ObHom}_C(a, b)$

$\text{isPrecategory } (C : \text{cat}) : U$

$= (\text{id} : (x : C.1) \rightarrow C.2 x x)$   
 $* (c : (x y z : C.1) \rightarrow C.2 x y \rightarrow C.2 y z \rightarrow C.2 x z)$   
 $* (\text{homSet} : (x y : C.1) \rightarrow \text{isSet } (C.2 x y))$   
 $* (\text{left} : (x y : C.1) \rightarrow (f : C.2 x y)$   
 $\rightarrow \text{Path } (C.2 x y) (c x x y (\text{id } x) f) f)$   
 $* (\text{right} : (x y : C.1) \rightarrow (f : C.2 x y)$   
 $\rightarrow \text{Path } (C.2 x y) (c x y y f (\text{id } y)) f)$   
 $* ((x y z w : C.1) (f : C.2 x y) (g : C.2 y z)$   
 $(h : C.2 z w) \rightarrow \text{Path } (C.2 x w)$   
 $(c x z w (c x y z f g) h) (c x y w f (c y z w g h)))$

## ObHom

```

carrier (C: precategory): U = C.1.1
hom      (C: precategory) (a b: carrier C): U = C.1.2 a b
path     (C: precategory) (x: carrier C): hom C x x = C.2.1 x
compose (C: precategory) (x y z: carrier C)
         (f: hom C x y) (g: hom C y z): hom C x z = C.2.2.1 x y z f g

```

$Ob_C \prod_{x,y:Ob_C} isContr(Hom_C(x,y))$

$Ob_C \prod_{x,y:Ob_C} isContr(Hom_C(y,x))$

```

isInitial (C: precategory) (x: carrier C): U
  = (y: carrier C) -> isContr (hom C x y)
isTerminal (C: precategory) (y: carrier C): U
  = (x: carrier C) -> isContr (hom C x y)
initial (C: precategory): U
  = (x: carrier C) * isInitial C x
terminal (C: precategory): U
  = (y: carrier C) * isTerminal C y

```

$ABF : A \rightarrow BF_{Ob} : Ob_H A \rightarrow Ob_B a, b : Ob_A F_{Hom} : Hom_A(a, b) \rightarrow$   
 $Hom_B(F_{Ob}(a), F_{Ob}(b)) a : Ob_A F_{Ob}(1_a) = 1_{F_{Ob}}(a), b, c : Ob_A f :$   
 $Hom_A(a, b) g : Hom_A(b, c) F(g \circ f) = F_{Hom}(g) \circ F_{Hom}(f)$

```

catfunctor (A B: precategory): U
  = (ob: carrier A -> carrier B)
  * (mor: (x y: carrier A) -> hom A x y -> hom B (ob x) (ob y))
  * (id: (x: carrier A) -> Path (hom B (ob x) (ob x))
      (mor x x (path A x)) (path B (ob x)))
  * ((x y z: carrier A) -> (f: hom A x y) -> (g: hom A y z) ->
      Path (hom B (ob x) (ob z)) (mor x z (compose A x y z f g))
      (compose B (ob x) (ob y) (ob z) (mor x y f) (mor y z g)))

```

$F, G : C \rightarrow D \gamma : F \rightarrow G x : C \gamma_a : Hom_D(F(x), G(x)) x, y : Cf : Hom_C(x, y)$   
 $G(f) \circ \gamma_x = \gamma_y \circ F(g)$

```

isNaturalTrans (C D: precategory)
  (F G: catfunctor C D)
  (eta: (x: carrier C) -> hom D (F.1 x) (G.1 x)): U
  = (x y: carrier C) (h: hom C x y) ->
    Path (hom D (F.1 x) (G.1 y))
      (compose D (F.1 x) (F.1 y) (G.1 y) (F.2.1 x y h) (eta y))
      (compose D (F.1 x) (G.1 x) (G.1 y) (eta x) (G.2.1 x y h))

```

```

ntrans (C D: precategory) (F G: catfunctor C D): U
  = (eta: (x: carrier C) -> hom D (F.1 x) (G.1 x))
  * (isNaturalTrans C D F G eta)

```

```

extension (C C' D: precategory)
  (K: catfunctor C C') (G: catfunctor C D) : U
  = (F: catfunctor C' D)
  * (ntrans C D (compFunctor C C' D K F) G)

```

$f : \text{Hom}_A(a, b) g : \text{Hom}_A(b, a) 1_a =_{\eta} g \circ f \circ g =_{\epsilon} 1_b = g a, b : A$   
 $a = b \rightarrow \text{iso}_A(a, b)$

```

iso (C: precategory) (A B: carrier C): U
  = (f: hom C A B)
  * (g: hom C B A)
  * (eta: Path (hom C A A) (compose C A B A f g) (path C A))
  * (Path (hom C B B) (compose C B A B g f) (path C B))

```

$a : \text{Ob}_C \prod_{A: \text{Ob}_C} \text{isContr}_{\Sigma_{B: \text{Ob}_C}} \text{iso}_C(A, B)$

```

isCategory (C: precategory): U
  = (A: carrier C) -> isContr ((B: carrier C) * iso C A B)
  category: U = (C: precategory) * isCategory C

```

```

Product (X Y: precategory) : precategory
Coproduct (X Y: precategory) : precategory

```

$CC^{\text{op}}$

```

opCat (P: precategory): precategory

```

```

sliceCat (C D: precategory)
  (a: carrier (opCat C))
  (F: catfunctor D (opCat C))
  : precategory
  = cosliceCat (opCat C) D a F

```

```

cosliceCat (C D: precategory)
(a: carrier C)
(F: catfunctor D C) : precategory

```

```

initArr (C D: precategory)
(a: carrier C)
(F: catfunctor D C): U = initial (cosliceCat C D a F)

```

```

termArr (C D: precategory)
(a: carrier (opCat C))
(F: catfunctor D (opCat C)): U = terminal (sliceCat C D a F)

```

$$\text{Ob} = \top \text{Hom} = \top$$

```

unitCat: precategory

```

```

Set: precategory = ((Ob,Hom),id,c,HomSet,L,R,Q) where
Ob: U = SET
Hom (A B: Ob): U = A.1 -> B.1
id (A: Ob): Hom A A = idfun A.1
c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
= o A.1 B.1 C.1 g f
HomSet (A B: Ob): isSet (Hom A B) = setFun A.1 B.1 B.2
L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f
= refl (Hom A B) f
R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f
= refl (Hom A B) f
Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
: Path (Hom A D) (c A C D (c A B C f g) h) (c A B D f (c B C D g h))
= refl (Hom A D) (c A B D f (c B C D g h))

```

```

Functions (X Y: U) (Z: isSet Y): precategory
= ((Ob,Hom),id,c,HomSet,L,R,Q) where
Ob: U = X -> Y
Hom (A B: Ob): U = id (X -> Y)
id (A: Ob): Hom A A = idfun (X -> Y)
c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C = idfun (X -> Y)

```

```

HomSet (A B: Ob): isSet (Hom A B) = setFun Ob Ob (setFun X Y Z)
L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f = axiom
R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f = axiom
Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
  : Path (Hom A D) (c A C D (c A B C f g) h)
    (c A B D f (c B C D g h)) = axiom

```

```

Cat: precategory = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = precategory
  Hom (A B: Ob): U = catfunctor A B
  id (A: Ob): catfunctor A A = idFunctor A
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
    = compFunctor A B C f g
  HomSet (A B: Ob): isSet (Hom A B) = axiom
  L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f = axiom
  R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f = axiom
  Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h)
      (c A B D f (c B C D g h)) = axiom

```

```

Func (X Y: precategory): precategory
  = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = catfunctor X Y
  Hom (A B: Ob): U = ntrans X Y A B
  id (A: Ob): ntrans X Y A A = axiom
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C = axiom
  HomSet (A B: Ob): isSet (Hom A B) = axiom
  L (A B: Ob) (f: Hom A B): Path (Hom A B) (c A A B (id A) f) f = axiom
  R (A B: Ob) (f: Hom A B): Path (Hom A B) (c A B B f (id B)) f = axiom
  Q (A B C D: Ob) (f: Hom A B) (g: Hom B C) (h: Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h)
      (c A B D f (c B C D g h)) = axiom

```

$k(k-1)01$

```

equiv: U
functor (C D: cat): U
ntrans (C D: cat) (F G: functor C D): U
modification (C D: cat) (F G: functor C D) (I J: ntrans C D F G): U

```

```

Cat2 : U
= (Ob: U)
* (Hom: (A B: Ob) -> U)
* (Hom2: (A B: Ob) -> (C F: Hom A B) -> U)
* (id: (A: Ob) -> Hom A A)
* (id2: (A: Ob) -> (B: Hom A A) -> Hom2 A A B B)
* (c: (A B C: Ob) (f: Hom A B) (g: Hom B C) -> Hom A C)
* (c2: (A B: Ob) (X Y Z: Hom A B)
  (f: Hom2 A B X Y) (g: Hom2 A B Y Z) -> Hom2 A B X Z)

```

$$A \xrightarrow{f} C \xleftarrow{g} BA \times_C Bpb_1 : \times_C \rightarrow Apb_2 : \times_C \rightarrow B$$

$$A \times_C B \xrightarrow{pb_2} B$$

$$A \xrightarrow{f} C$$

$(\times_C, pb_1, pb_2)(D, q_1, q_2)u : D \rightarrow \times_C pb_1 \circ u = q_1 pb_2 \circ q_2$

```

homTo (C: precategory) (X: carrier C): U
= (Y: carrier C) * hom C Y X
cospan (C: precategory): U
= (X: carrier C) * ( _: homTo C X) * homTo C X
cospanCone (C: precategory) (D: cospan C): U
= (W: carrier C) * hasCospanCone C D W
cospanConeHom (C: precategory) (D: cospan C)
  (E1 E2: cospanCone C D) : U
= (h: hom C E1.1 E2.1) * isCospanConeHom C D E1 E2 h
isPullback (C: precategory) (D: cospan C) (E: cospanCone C D) : U
= (h: cospanCone C D) -> isContr (cospanConeHom C D h E)
hasPullback (C: precategory) (D: cospan C) : U
= (E: cospanCone C D) * isPullback C D E

```

$ABF : A \rightarrow BF_{Ob} : Ob_h A \rightarrow Ob_B a, b : Ob_A F_{Hom} : Hom_A(a, b) \rightarrow$   
 $Hom_B(F_{Ob}(a), F_{Ob}(b))a : Ob_A F_{Ob}(1_a) = 1_{F_{Ob}}(a)a, b, c : Ob_A f :$   
 $Hom_A(a, b)g : Hom_A(b, c)F(g \circ f) = F_{Hom}(g) \circ F_{Hom}(f)$

```

catfunctor (A B: precategory): U
= (ob: carrier A -> carrier B)
* (mor: (x y: carrier A) -> hom A x y -> hom B (ob x) (ob y))
* (id: (x: carrier A) -> Path (hom B (ob x) (ob x))
  (mor x x (path A x)) (path B (ob x)))

```

```

* ((x y z: carrier A) -> (f: hom A x y) -> (g: hom A y z) ->
  Path (hom B (ob x) (ob z)) (mor x z (compose A x y z f g))
    (compose B (ob x) (ob y) (ob z) (mor x y f) (mor y z g)))

```

$\text{Ob}_C$

$$\prod_{x,y:\text{Ob}_C} \text{isContr}(\text{Hom}_C(y,x)).$$

```

isTerminal (C: precategory) (y: carrier C): U
= (x: carrier C) -> isContr (hom C x y)
terminal (C: precategory): U
= (y: carrier C) * isTerminal C y

```

$\infty$

$\text{PROP}_{x,y} : P_{x=y}$

$$\text{isProp}(P) = \prod_{x,y:P} (x = y).$$

$x, y : A, p, q : x =_A y \Rightarrow p = q$

$x, y : A, p, q : x =_A y, r, s : p =_{=A} q \Rightarrow r = s$

$\text{SET}_{\text{SET}} x, y : A, p, q : x = y \Rightarrow p = q$

$$\text{isSet}(A) = \prod_{x,y:A} \prod_{p,q:x=y} (p = q).$$

`data N = Z | S (n: N)`

```

n_grpd (A: U) (n: N): U = (a b: A) -> rec A a b n where
  rec (A: U) (a b: A) : (k: N) -> U
    = split { Z -> Path A a b ; S n -> n_grpd (Path A a b) n }

```

```

isContr (A: U): U = (x: A) * ((y: A) -> Path A x y)
isProp (A: U): U = n_grpd A Z
isSet (A: U): U = n_grpd A (S Z)
PROP : U = (X:U) * isProp X
SET : U = (X:U) * isSet X

```

$\Pi$

```

setPi (A: U) (B: A -> U) (h: (x: A) -> isSet (B x)) (f g: Pi A B)
  (p q: Path (Pi A B) f g)
  : Path (Path (Pi A B) f g) p q

```

$\Sigma\Sigma$

```

setSig (A:U) (B: A -> U) (base: isSet A)
  (fiber: (x:A) -> isSet (B x)) : isSet (Sigma A B)

```

```

data unit = tt
unitRec (C: U) (x: C): unit -> C = split tt -> x
unitInd (C: unit -> U) (x: C tt): (z:unit) -> C z
  = split tt -> x

```

## SetHom $\Pi$

```

Set: precategory = ((Ob,Hom),id,c,HomSet,L,R,Q) where
  Ob: U = SET
  Hom (A B: Ob): U = A.1 -> B.1
  id (A: Ob): Hom A A = idfun A.1
  c (A B C: Ob) (f: Hom A B) (g: Hom B C): Hom A C
    = o A.1 B.1 C.1 g f
  HomSet (A B: Ob): isSet (Hom A B) = setFun A.1 B.1 B.2
  L (A B:Ob) (f:Hom A B): Path (Hom A B)(c A A B (id A)f)f
    = refl (Hom A B) f
  R (A B:Ob) (f:Hom A B): Path (Hom A B)(c A B B f(id B))f
    = refl (Hom A B) f
  Q (A B C D: Ob) (f:Hom A B) (g:Hom B C) (h:Hom C D)
    : Path (Hom A D) (c A C D (c A B C f g) h)
      (c A B D f (c B C D g h))
    = refl (Hom A D) (c A B D f (c B C D g h))

```

## AS $\in$ ASSSSSS

```

Structure topology (A : Type) := {
  open :> (A -> Prop) -> Prop;
  empty_open: open (empty _);
  full_open: open (full _);
  inter_open: forall u,
    open u -> forall v, open v
      -> open (inter A u v) ;
  union_open: forall s, (subset _ s open)
    -> open (union A s) }.

```

$$R \subset \text{Hom}_{\mathcal{C}}(\cdot, U), U \in \mathcal{C},$$

$$R \subset \text{Hom}_{\mathcal{C}}(\cdot, U)\phi : V \rightarrow UC$$

$$\phi^{-1}(R) = \{\gamma : W \rightarrow V \mid \phi \cdot \gamma \in R\}$$

$$VR, R' \subset \text{Hom}_{\mathcal{C}}(\cdot, U)\phi^{-1}(R')\phi : V \rightarrow URR' \text{Hom}_{\mathcal{C}}(\cdot, U)U \in \mathcal{C}$$

$$\text{Ob}_{\mathcal{C}}\{f_i : U_i \rightarrow U\}_{i \in I}g : V \rightarrow U\{h : V_j \rightarrow V\}_{j \in J}h_j \circ g f_i \frac{V_j U_i}{V \rightarrow U}$$

Co (C: **precategory**) (cod: carrier C) : U  
 = (dom: carrier C)  
 \* (hom C dom cod)

Delta (C: **precategory**) (d: carrier C) : U  
 = (index: U)  
 \* (index -> Co C d)

Coverage (C: **precategory**): U  
 = (cod: carrier C)  
 \* (fam: Delta C cod)  
 \* (coverings: carrier C -> Delta C cod -> U)  
 \* (coverings cod fam)

CCC

$$\{\phi_{\alpha} : U_{\alpha} \rightarrow U\}, U \in \mathcal{C},$$

$$\phi_{\alpha} : U_{\alpha} \rightarrow U \psi : V \rightarrow UC \times_U U_{\alpha} \rightarrow VV\{\phi_{\alpha} : U_{\alpha} \rightarrow U\}\{\gamma_{\alpha, \beta} : W_{\alpha, \beta} \rightarrow U_{\alpha}\}\alpha$$

$$W_{\alpha, \beta} \xrightarrow{\gamma_{\alpha, \beta}} U_{\alpha} \xrightarrow{\phi_{\alpha}} U$$

$$\{1 : U \rightarrow U\}U \in \mathcal{C}$$

site (C: **precategory**): U  
 = (C: **precategory**) \* Coverage C

$$CC^{op} \rightarrow \text{Set}$$

presheaf (C: **precategory**): U  
 = **catfunctor** (opCat C) Set

C

$$F : C^{op} \rightarrow \text{Set}$$

$$F(U) \rightarrow \lim_{V \rightarrow U \in R}^{\leftarrow} F(V)$$

$$R \subset \text{Hom}_{\mathcal{C}}(\cdot, U)$$

$$\text{Hom}_{\mathcal{C}}(\text{Hom}_{\mathcal{C}}(\cdot, U), F) \rightarrow \text{Hom}_{\mathcal{C}}(R, F)$$

```

sheaf (C: precategory): U
= (S: site C)
* presheaf S.1

```

$(C, J)C$

$CR \rightarrow ER \rightarrow E \rightarrow QC$

CD

$f : (C) \rightarrow (D)$

$f_* : (C) \rightarrow (D) f^* : (D) \rightarrow (C) f^* f_* f^* f_* f^*$

$ES(p^*, p_*) : E \rightarrow Sp^1 \vdash p_* p^1 \dashv p_* p^* p^1 p_!$

$\int \dashv b \dashv \#$

$f : Y \rightarrow ZX g_1, g_2 : X \rightarrow Y$

$f \circ g_1 = f \circ g_2 \rightarrow g_1 = g_2.$

$XHom(X, )Hom(X, Y) \rightarrow Hom(X, Z)$

```

mono (P: precategory) (Y Z: carrier P) (f: hom P Y Z): U
= (X: carrier P) (g1 g2: hom P X Y)
-> Path (hom P X Z) (compose P X Y Z g1 f)
        (compose P X Y Z g2 f)
-> Path (hom P X Y) g1 g2

```

$Ctrue : 1 \rightarrow \Omega 1U \rightarrow X_{XU} : X \rightarrow \Omega_{XU} \overset{k_1}{\dashv} \overset{k_2}{\dashv} \overset{k_3}{\dashv} \overset{k_4}{\dashv}$

```

subobjectClassifier (C: precategory): U
= (omega: carrier C)
* (end: terminal C)
* (trueHom: hom C end.1 omega)
* (chi: (V X: carrier C) (j: hom C V X) -> hom C X omega)
* (square: (V X: carrier C) (j: hom C V X) -> mono C V X j
-> hasPullback C (omega, (end.1, trueHom), (X, chi V X j)))
* ((V X: carrier C) (j: hom C V X) (k: hom C X omega)
-> mono C V X j
-> hasPullback C (omega, (end.1, trueHom), (X, k))
-> Path (hom C X omega) (chi V X j) k)

```

CMLTT

```

isCCC (C: precategory): U
= (Exp: (A B: carrier C) -> carrier C)
* (Prod: (A B: carrier C) -> carrier C)
* (Apply: (A B: carrier C) -> hom C (Prod (Exp A B) A) B)
* (P1: (A B: carrier C) -> hom C (Prod A B) A)
* (P2: (A B: carrier C) -> hom C (Prod A B) B)
* (Term: terminal C)
* unit

```

## $\Pi\Sigma$ SETisSetpropPiMLTT

```

cartesianClosure : isCCC Set
= (expo,prod,appli,proj1,proj2,term,tt) where
  exp (A B: SET): SET = (A.1 -> B.1, setFun A.1 B.1 B.2)
  pro (A B: SET): SET = (prod A.1 B.1, setSig A.1 (\(_ : A.1)
    -> B.1) A.2 (\(_ : A.1) -> B.2))
  expo: (A B: SET) -> SET = \(\A B: SET) -> exp A B
  prod: (A B: SET) -> SET = \(\A B: SET) -> pro A B
  appli: (A B: SET) -> hom Set (pro (exp A B) A) B
    = \(\A B: SET) -> \(\x:(pro (exp A B) A).1)-> x.1 x.2
  proj1: (A B: SET) -> hom Set (pro A B) A
    = \(\A B: SET) (x: (pro A B).1) -> x.1
  proj2: (A B: SET) -> hom Set (pro A B) B
    = \(\A B: SET) (x: (pro A B).1) -> x.2
  unitContr (x: SET) (f: x.1 -> unit) : isContr (x.1 -> unit)
    = (f, \(\z: x.1 -> unit) -> propPi x.1 (\(\_:x.1)->unit)
      (\(\x:x.1) -> propUnit) f z)
  term: terminal Set = ((unit,setUnit),
    \(\x: SET) -> unitContr x (\(\z: x.1) -> tt))

```

```

Topos (cat: precategory) : U
= (cartesianClosure: isCCC cat)
* subobjectClassifier cat

```

```

internal : Topos Set
= (cartesianClosure,hasSubobject)

```

$$Cb : \text{Ob}_C \downarrow bf : a \rightarrow bf^* : C \downarrow b \rightarrow c \downarrow a \sum_f \prod_f$$

$$S^{n-1} \hookrightarrow D^n \hookrightarrow D^n \hookrightarrow S^{n-1}$$

$$X_f : S^{n-1} \rightarrow X$$

$$\begin{array}{c} S^{n-1} \xrightarrow{X_f} X \\ \uparrow \\ D^n \xrightarrow{f} D^n \end{array}$$

$$X \cup_f D^n$$

$$-1 \leq i \leq n-1$$

$$\text{Xcolimit}(X_i)_{X_{-1} = \emptyset} \hookrightarrow X_0 \hookrightarrow X_1 \hookrightarrow X_2 \hookrightarrow \dots \hookrightarrow X_i \hookrightarrow X_{i+1} \hookrightarrow X_i$$

$$\emptyset \hookrightarrow X_0 \hookrightarrow X_1 \hookrightarrow X_2 \hookrightarrow \dots \hookrightarrow X$$

$$X_i \leq i$$

$$(A, a)A : Ua : A$$

$$\text{pointed} : U = (A : U) * A$$

$$\text{point} (A : \text{pointed}) : A.1 = A.2$$

$$\text{space} (A : \text{pointed}) : U = A.1$$

$$\Omega(A, a) =_{\text{def}} ((a =_A a), \text{refl}_A(a)).$$

$$\text{omega1} (A : \text{pointed}) : \text{pointed}$$

$$= (\text{Path} (\text{space } A) (\text{point } A) (\text{point } A), \text{refl } A.1 (\text{point } A))$$

$$\begin{cases} \Omega^0(A, a) =_{\text{def}} (A, a) \\ \Omega^{n+1}(A, a) =_{\text{def}} \Omega^n(\Omega(A, a)) \end{cases}$$

$$\text{omega} : \text{nat} \rightarrow \text{pointed} \rightarrow \text{pointed} = \text{split}$$

$$\text{zero} \rightarrow \text{idfun } \text{pointed}$$

$$\text{succ } n \rightarrow \lambda(A : \text{pointed}) \rightarrow \text{omega } n (\text{omega1 } A)$$

$$\pi_n S^m = \|\Omega^n(S^m)\|_0.$$

```

piS (n: nat): (m: nat) -> U = split
  zero  -> sTrunc (space (omega n (bool,false)))
  succ x -> sTrunc (space (omega n (Sn (succ x),north)))

```

$$\Omega(S^1) = \mathbb{Z}$$

```

data S1 = base
  | loop <i> [ (i=0) -> base ,
              (i=1) -> base ]

loopS1 : U = Path S1 base base

encode (x:S1) (p:Path S1 base x)
  : helix x
  = subst S1 helix base x p zeroZ

decode : (x:S1) -> helix x -> Path S1 base x = split
  base -> loopIt
  loop @ i -> rem @ i where
    p : Path U (Z -> loopS1) (Z -> loopS1)
    = <j> helix (loop1@j) -> Path S1 base (loop1@j)
  rem : PathP p loopIt loopIt
    = corFib1 S1 helix (\(x:S1)->Path S1 base x) base
      loopIt loopIt loop1 (\(n:Z) ->
        comp (<i> Path loopS1 (oneTurn (loopIt n))
              (loopIt (testIsoPath Z Z sucZ predZ
                          sucPredZ predsucZ n @ i)))
              (<i>(lem1It n)@-i) [])

loopS1eqZ : Path U Z loopS1
  = isoPath Z loopS1 (decode base) (encode base)
  sectionZ retractZ

```

$$S^3 S^1 S^2 S^3 R^3 S^0 S^1 S^3 S^7$$

$$S^3 S^2$$

$$S^3$$

$$S^3 \mathbb{R}^4$$

$$S^3 = \{(x_0, x_1, x_2, x_3) \in \mathbb{R}^4 : \sum_{i=0}^3 x_i^2 = 1\};$$

$$\mathbb{H}$$

$$S^3 = \{x \in \mathbb{H} : \|x\| = 1\}.$$

$S^3(\eta, \theta_1, \theta_2)$

$$\begin{cases} x_0 = \cos(\theta_1)\sin(\eta), \\ x_1 = \sin(\theta_1)\sin(\eta), \\ x_2 = \cos(\theta_2)\cos(\eta), \\ x_3 = \sin(\theta_2)\cos(\eta). \end{cases}$$

$\eta \in [0, \frac{\pi}{2}] \theta_{1,2} \in [0, 2\pi]$

$S^2 S^2 \theta_2$

$$\begin{cases} x = \sin(2\eta)\cos(\theta_1), \\ y = \sin(2\eta)\sin(\theta_1), \\ z = \cos(2\eta). \end{cases}$$

```
var fiber = new THREE.Curve(),
    color = sphericalCoords.color;

fiber.getPoint = function(t) {
  var eta = sphericalCoords.eta,
      phi = sphericalCoords.phi,
      theta = 2 * Math.PI * t;
  var x1 = Math.cos(phi+theta) * Math.sin(eta/2),
      x2 = Math.sin(phi+theta) * Math.sin(eta/2),
      x3 = Math.cos(phi-theta) * Math.cos(eta/2),
      x4 = Math.sin(phi-theta) * Math.cos(eta/2);
  var m = mag([x1,x2,x3]),
      r = Math.sqrt((1-x4)/(1+x4));
  return new THREE.Vector3(r*x1/m,r*x2/m, r*x3/m);
};
```

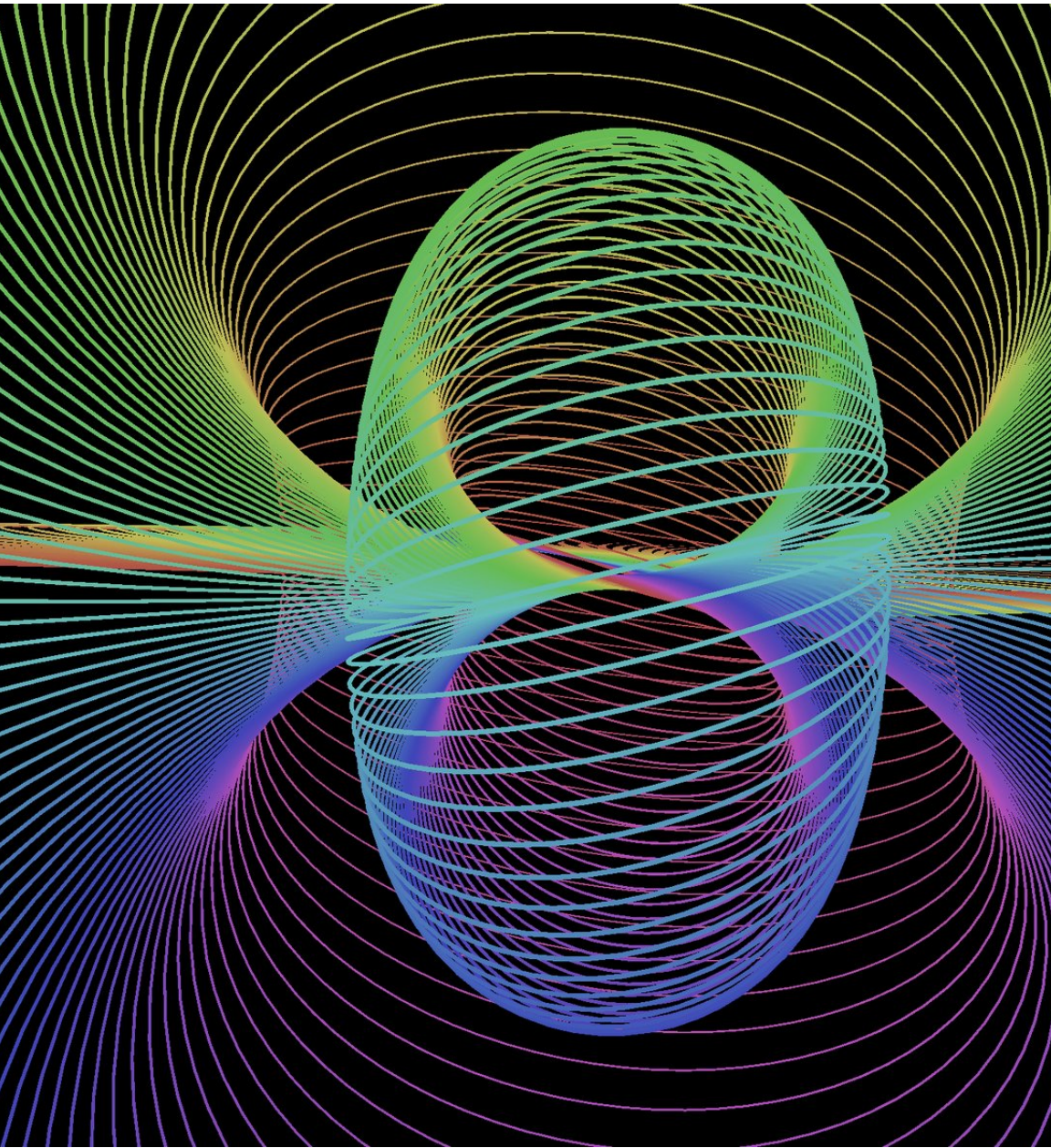
$S^3$

```
rot: (x : S1) -> Path S1 x x = split
  base -> loop1
  loop @ i -> constSquare S1 base loop1 @ i

mu : S1 -> equiv S1 S1 = split
  base -> idEquiv S1
  loop @ i -> equivPath S1 S1 (idEquiv S1)
    (idEquiv S1) (<j> \ (x : S1) -> rot x @ j) @ i

H : S2 -> U = split
  north -> S1
  south -> S1
  merid x @ i -> ua S1 S1 (mu x) @ i

total : U = (c : S2) * H c
```



A

$$H_A = \begin{cases} A : \mathbb{U} \\ e : A \\ \mu : A \rightarrow A \rightarrow A \\ \beta : (a : A) \rightarrow \Sigma(\mu(e, a) = a)(\mu(a, e) = a) \end{cases}$$

$$(S^0, S^1, p, S^1)(S^1, S^3, p, S^2)(S^3, S^7, p, S^4)(S^7, S^{15}, p, S^8)$$

$$\phi : S^{2^n-1} \rightarrow S^n \text{ cofib}(\phi) = S^n \cup_{\phi} \mathbb{D}^{2^n}$$

$$H^k(\text{cofib}(\phi), \mathbb{Z}) = \begin{cases} \mathbb{Z} & \text{for } k = n, 2n \\ 0 & \text{otherwise} \end{cases}$$

$$\alpha, \beta \in \pi_n \text{ cofib}(\phi) \quad \alpha \beta \alpha \sqcup \alpha = h(\phi) \cdot \beta h(\phi)$$

1

---

## O<sub>HTS</sub>

```
module buddhism where
import path

concept (o: U): U
  = o -> U

nondual (o: U) (p: concept o): U
  = (x y: o) -> Path U (p x) (p y)

allpaths (o: U): U
  = (x y: o) -> Path o x y

  =U=

encode (o:U): ((p: concept o) -> nondual o p) -> allpaths o
  = \nd: (p: concept o) -> nondual o p) (a b: o)
  -> coerce(Path o a a)(Path o a b)(nd(\(z:o)->Path o a z)a b)(refl o a)

decode (o:U): allpaths o -> ((p: concept o) -> nondual o p)
  = \(all: allpaths o)(p: concept o)(x y: o) -> cong o U p x y (all x y)

1
```

```

CoInductive Co (E : Effect.t) : Type -> Type :=
| Bind : forall (A B : Type), Co E A -> (A -> Co E B) -> Co E B
| Split : forall (A : Type), Co E A -> Co E A -> Co E A
| Join : forall (A B : Type), Co E A -> Co E B -> Co E (A * B).
| Ret : forall (A : Type) (x : A), Co E A
| Call : forall (command : Effect.command E),
          Co E (Effect.answer E command)

Definition run (argv : list LString.t): Co effect unit :=
ido! log (LString.s "What is your name?") in
ilet! name := read_line in
match name with
| None => ret tt
| Some name => log (LString.s "Hello " ++ name ++ LString.s "!")
end.

Parameter infinity : nat.
Definition eval {A} (x : Co effect A) : Lwt.t A := eval_aux infinity x.

Fixpoint eval_aux {A} (s: nat) (x: Co effect A) : Lwt.t A :=
match s with
| 0 => error tt
| S s =>
match x with
| Bind _ _ x f => Lwt.bind (eval_aux s x) (fun v_x => eval_aux s (f
v_x))
| Split _ x y => Lwt.choose (eval_aux s x) (eval_aux s y)
| Join _ _ x y => Lwt.join (eval_aux s x) (eval_aux s y)
| Ret _ v => Lwt.ret v
| Call c => eval_command c
end
end.

CoFixpoint handle_commands : Co effect unit :=
ilet! name := read_line in
match name with
| None => ret tt
| Some command =>
ilet! result := log (LString.s "Input: "
++ command ++ LString.s ".")
in handle_commands
end.

Definition launch (m: list LString.t -> Co effect unit): unit :=
let argv := List.map String.to_lstring Sys.argv in
Lwt.launch (eval (m argv)).

Definition corun (argv: list LString.t): Co effect unit :=
handle_commands.

Definition main := launch corun.

```

```

String: Type = List Nat
data IO: Type =
  (getLine: (String -> IO) -> IO)
  (putLine: String -> IO)
  (pure: () -> IO)

-- IOI/@: (r: U) [x: U] [[s: U] -> s -> [s -> #IOI/F r s] -> x] x
  \ (r : *)
-> \ (x : *)
-> (\ (s : *)
  -> s
  -> (s -> #IOI/F r s)
  -> x)
-> x

-- IOI/F
  \ (a : *)
-> \ (State : *)
-> \ (IOF : *)
-> \ (PutLine_ : #IOI/data -> State -> IOF)
-> \ (GetLine_ : (#IOI/data -> State) -> IOF)
-> \ (Pure_ : a -> IOF)
-> IOF

-- IOI/MkIO
  \ (r : *)
-> \ (s : *)
-> \ (seed : s)
-> \ (step : s -> #IOI/F r s)
-> \ (x : *)
-> \ (k : forall (s : *) -> s -> (s -> #IOI/F r s) -> x)
-> k s seed step

-- Morte/corecursive
( \ (r: *1)
-> ( (((#IOI/MkIO r) (#Maybe/@ #IOI/data)) (#Maybe/Nothing #IOI/data))
  ( \ (m: (#Maybe/@ #IOI/data))
  -> ((((#Maybe/maybe #IOI/data) m) ((#IOI/F r) (#Maybe/@ #IOI/data)))
    ( \ (str: #IOI/data)
    -> (((#IOI/putLine r) (#Maybe/@ #IOI/data)) str)
      (#Maybe/Nothing #IOI/data))))
  (((#IOI/getLine r) (#Maybe/@ #IOI/data))
  (#Maybe/Just #IOI/data))))))

```

```

copure() ->
  fun (_) -> fun (IO) -> IO end end.

cogetLine() ->
  fun (IO) -> fun (_) ->
    L = ch:list(io:get_line("> ")),
    ch:ap(IO, [L]) end end.

coputLine() ->
  fun (S) -> fun (IO) ->
    X = ch:unlist(S),
    io:put_chars(" ++X),
    case X of "0\n" -> list([]);
      _ -> corec() end end end.

corec() ->
  ap('Morte':corecursive(),
    [copure(), cogetLine(), coputLine(), copure(), list([])]).

```

```

> om_extract:extract("priv/normal/IOI").
ok
> Active: module loaded: {reloaded,'IOI'}

```

```

> om:corec().
> 1
: 1
> 0
: 0
#Fun<List.3.113171260>

```

```

-- IO/@
  \ (a : *)
-> \ / (IO : *)
-> \ / (GetLine_ : (#IO/data -> IO) -> IO)
-> \ / (PutLine_ : #IO/data -> IO -> IO)
-> \ / (Pure_ : a -> IO)
-> IO

```

```

-- IO/replicateM
  \ (n: #Nat/@)
-> \ (io: #IO/@ #Unit/@)
-> #Nat/fold n (#IO/@ #Unit/@)
    (#IO/[>>] io)
    (#IO/pure #Unit/@ #Unit/Make)

```

```

-- Morte/recursive
((#IO/replicateM #Nat/Five)
  (((#IO/[>>] #IO/data) #Unit/@) #IO/getLine) #IO/putLine))

```

```
pure() ->
  fun(IO) -> IO end.

getLine() ->
  fun(IO) -> fun(_) ->
    L = ch:list(io:get_line("> ")),
    ch:ap(IO, [L]) end end.

putLine() ->
  fun (S) -> fun(IO) ->
    io:put_chars(" ++ch:unlist(S)),
    ch:ap(IO, [S]) end end.

rec() ->
  ap('Morte':recursive(),
    [getLine(),putLine(),pure(),list([])]).

> om:rec().
> 1
: 1
> 2
: 2
> 3
: 3
> 4
: 4
> 5
: 5
#Fun<List.28.113171260>
```



